Programming Languages and Techniques (CIS120)

Lecture 4

Lists, Recursion and Tuples

Announcements

- Read Chapters 3 (Lists) and 4 (Tuples) of the lecture notes
- We will start Chapters 5 & 6 on Monday
- HW01 due Tuesday at midnight
- Poll Everywhere attendance starts Monday

Review: What is a list?

- The '::' operator, pronounced "cons", *constructs* a new list from a head element and a shorter list.
- Lists are an example of an *inductive datatype*.

Calculating with Matches

• Consider how to evaluate a match expression: begin match [1;2;3] with $|$ $|$ \rightarrow 42 | first:: rest \rightarrow first $+10$ end

The Inductive Nature of Lists

- What is going on!? The definition of list mentions 'list'!
- Insight: 'list' is *inductive*:
	- The empty list [] is the (only) list of 0 elements
	- To construct a list of (1+n) elements, add an element to an *existing* list of n elements
	- The set of list values contains *all and only* values constructed this way
- Corresponding computation principle: *recursion*

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Recursion

Recursion principle:

Compute a function value for a given input by combining the results for strictly smaller subcomponents of the input.

- The structure of the computation follows the inductive structure of the input.
- Example:
	-
	-
	-
	- l ength \Box = 0
- $length (1::2::3::[]) = 1 + (length (2::3::[]))$
- $length (2::3::[]) = 1 + (length (3::[]))$
- $length (3::[])$ = 1 + (length $[])$
	-

Recursion Over Lists

If the list is non-empty, then "x" is the first string in the list and "rest" is the remainder of the list.

Calculating with Recursion

```
length \lceil "a"; "b"]
```

```
⟼ (substitute the list for l in the function body)
        begin match "a"::"b"::<sup>[]</sup> with
        | \Box -> 0
        | ( x :: rest ) -> 1 + length rest
        end
```

```
⟼ (second case matches with rest = "b"::[])
```
 $1 + (length "b":[])$

```
⟼ (substitute the list for l in the function body)
```

```
1 + (begin match "b":: \Box with
        | | | \rightarrow 0
        | ( x :: rest ) -> 1 + length rest
       end )
```
⟼ *(second case matches again, with rest = [])*

```
1 + (1 + length)
```

```
⟼ (substitute [] for l in the function body)
```

```
\mapsto 1 + 1 + 0 \Rightarrow 2
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```
…

```
let rec length (l:string list) : int=
  begin match l with
  | [] -> 0
  | ( x :: rest ) -> 1 + length rest
  end
```
More examples…

```
let rec sum (l : int list) : int =
  begin match l with
  | | | \rightarrow 0
  | ( x :: rest ) -> x + sum rest
  end
```

```
let rec contains (l:string list) (s:string):bool =
  begin match l with
  |\bigcap -> false| ( x :: rest ) \rightarrow s = x || contains rest s
  end
```
What best describes the behavior of (foo 3 l) ? It returns true if...

let rec foo (z:int) (l : int list): bool = begin match I with \Box -> true $|(x::rest") \rightarrow$ $(x > z)$ && foo z rest end

Y

1. Every element of lis less than 3.

2. Every element of lis greater than 3

3. There exists an element in I that is less than 3

4. There exists an element in I that is greater than 3

r.

What best describes the behavior of the function call (foo 3 l)?

ANSWER: every element is greater than 3

Structural Recursion Over Lists

Structural recursion builds an answer from smaller components:

```
let rec f (l : … list) … : … =
 begin match l with
  \Box -> ... (* BASE CASE *)
  | ( hd :: rest ) ->
      … (f rest) …. (* INDUCTIVE CASE *)end
```
The branch for $[$] calculates the value $(f [$ $])$ directly.

– this is the *base case* of the recursion

The branch for hd:: rest calculates $(f(hd::rest))$ given hd and $(f \text{ rest})$. – this is the *inductive case* of the recursion

Design Pattern for Recursion

1. Understand the problem What are the relevant concepts and how do they relate? 2. Formalize the interface How should the program interact with its environment? 3. Write test cases If the main input to the program is an immutable list, make sure the tests cover both empty and non-empty cases 4. Implement the required behavior If the main input to the program is an immutable list, look for a recursive solution… Is there a direct solution for the empty list? • Suppose someone has given us a partial solution that works for lists up to a certain size. Can we use it to build a better solution that works for lists that are one element larger?

Tuples and Tuple Patterns

Two Forms of Structured Data

OCaml provides two basic ways of packaging multiple values together into a single compound value:

- **Lists:**
	- *arbitrary-length* sequence of values of a *single type*
	- example: a list of email addresses
- **Tuples:**
	- *fixed-length* sequence of values, possibly of *different types*
	- example: tuple of name, phone #, and email

Tuples

• In OCaml, tuples are created by writing a sequence of expressions, separated by commas, inside parens:

```
let my pair = (3, true)let my triple = ('Hello'', 5, false)let my quadruple = (1, 2, "three", false)
```
• Tuple types are written using '*'

– e.g. my_triple has type:

string * int * bool

Cartesian Products

• The values of a tuple (a.k.a. product) type are tuples of elements from each component type.

Ocaml notation: $A * R$

Pattern Matching on Tuples

• Tuples can be inspected by pattern matching:

```
let first (x: string * int): string =
  begin match x with
  | (left, right) -> left
  end
first ("b", 10)
⇒
``h"
```
• As with lists, tuple patterns follow the syntax of tuple values and give names to the subcomponents so they can be used on the right-hand side of the ->

Mixing Tuples and Lists

• Tuples and lists can mix freely:

$$
[(1,"a''); (2,"b''); (3,"c'')]
$$

;
(int * string) list

$$
([1;2;3], ["a"; "b"; "c"])
$$

:: (int list) * (string list)

Nested Patterns

• We're seen several kinds of *simple patterns*:

• We can build *nested patterns* out of simple ones:

Wildcard Pattern

- Another handy simple pattern is the wildcard
- A wildcard pattern indicates that the value of the corresponding subcomponent is not used on the right-hand side of the match case.
	- And hence needs no name
	- _::tl *matches a non-empty list, but only names tail* (_,x) *matches a pair, but only names the 2nd part*

Unused Branches

- The branches in a match expression are considered in order from top to bottom.
- If you have "redundant" matches, then some later branches might not be reachable.
	- OCaml will give you a warning in this case

What is the value of this expression?

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Answer: 3

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Exhaustive Matches

- Pattern matching is *exhaustive* if there is a pattern for every possible value
- Example of a *non-exhaustive* match:

```
let sum_two (l : int list) : int =
  begin match l with
  | x::y::_ - \rightarrow x+yend
```
• OCaml will give you a warning and show an example of what isn't covered by your cases

Exhaustive Matches

- Pattern matching is *exhaustive* if there is a pattern for every possible value
- Example of an *exhaustive* match:

```
let sum_two (l : int list) : int =
  begin match l with
  | x::y::_ - \rightarrow x+y| _ -> failwith "not a length 2 list"
  end
```
• The wildcard pattern and failwith are useful tools for ensuring match coverage

More List & Tuple Programming

see patterns.ml

Example: zip

• zip takes two lists of the same length and returns a single list of pairs:

> zip $[1; 2; 3]$ $["a"; "b"; "c"] \Rightarrow$ $[(1, "a"); (2, "b"); (3, "c")]$

```
let rec zip (l1: int list)
             (l2: string list) : (int * string) list =
  begin match (l1, l2) with
  | (1, 1) \rightarrow 1(x:: xs, y:: ys) \rightarrow (x, y): (zip xs ys)| _ -> failwith "zip: unequal length lists"
  end
```