

# Programming Languages and Techniques (CIS120)

## Lecture 12

Partiality, Sequencing  
Chapters 11, 12

# Midterm 1

- Friday, September 27<sup>th</sup>
- Coverage: up to Monday, Sept. 23 (Chs. 1-10)
- Time: During lecture (001 @ 11am, 002 @ noon)  
Last names: A – L Leidy Labs 10  
Last names: M – Z Stitler (STIT) B6
- Review Session: TONIGHT 6:00-8:00pm Towne 100
- Review Material:
  - old exams on the web site lecture schedule
- Makeup exam
  - Monday, Sept. 30<sup>th</sup>
  - sign up form on the web site

# Announcements

- Dr. Sheth will have extra office hours  
Thursday 4:00-6:00PM in Levine 264
- Homework 4
  - Available soon after exam
  - Due: Tuesday, Oct. 8th

# Signature: Finite Map

```
module type MAP = sig

  type ('k, 'v) map

  val empty      : ('k, 'v) map
  val add        : 'k -> 'v -> ('k, 'v) map -> ('k, 'v) map
  val remove     : 'k          -> ('k, 'v) map -> ('k, 'v) map
  val mem        : 'k -> ('k, 'v) map -> bool
  val get        : 'k -> ('k, 'v) map -> 'v
  val entries    : ('k, 'v) map -> ('k * 'v) list
  val equals     : ('k, 'v) map -> ('k, 'v) map -> bool

end
```

# Properties of Finite Maps

For any finite map  $m$ , key  $k$ , and value  $v$ :

1.  $\text{get } k \text{ (add } k \ v \ m) = v$
2. If  $k_1 \neq k_2$  and  $\text{get } k_1 \ m = v_1$  then  
 $\text{get } k_1 \text{ (add } k_2 \ v_2 \ m) = v_1$
3. if  $\text{mem } k \ m = \text{true}$  then  
there is a  $v$  such that  $\text{get } k \ m = v$
4. If  $\text{mem } k \ m = \text{false}$  then  
 $\text{get } k \ m = v$  fails
5.  $\text{mem } k \text{ (add } k \ v \ m) = \text{true}$
6.  $\text{mem } k \text{ (remove } k \ m) = \text{false}$

And others...

# Completing module implementation

finiteMap.ml

# Implementation: Ordered Lists

```
module Assoc : MAP = struct
```

```
  (* Represent a finite map as a list of pairs. *)
  (* Representation invariant: *)
  (*   - no duplicate keys (helps get, remove) *)
  (*   - keys are sorted (helps equals, helps get) *)
```

```
type ('k, 'v) map = ('k * 'v) list
```

```
let empty : ('k, 'v) map = []
```

```
let rec mem (key:'k) (m : ('k, 'v) map) : bool =
  begin match m with
  | [] -> false
  | (k,v)::rest ->
      (key >= k) &&
      ((key = k) || (mem key rest))
  end
```

```
:: run_test "mem test" (fun () -> mem "b" [("a",3); ("b",4)])
```

# Implementation: Ordered Lists

```
let rec get (key:'k) (m : ('k,'v) map) : 'v =  
  begin match m with  
  | [] -> failwith "key not found"  
  | (k,v)::rest ->  
    if key < k then failwith "key not found"  
    else if key = k then v  
    else get key rest  
  end
```

```
let rec remove (key:'k) (m : ('k,'v) map) : ('k,'v) map =  
  begin match m with  
  | [] -> []  
  | (k,v)::rest ->  
    if key < k then m  
    else if key = k then rest  
    else (k,v)::remove key rest  
  end
```



# Abstract types

BIG IDEA: Hide the *concrete representation* of a type behind an *abstract interface* to preserve invariants

- The interface **restricts** how other parts of the program can interact with the data
  - Type checking ensures that the **only** way to create a set is with the operations in the interface
  - If all operations preserve invariants, then all sets in the program must satisfy invariants
  - Example: all BST-implemented sets must satisfy the BST invariant, therefore the lookup function can assume that its input satisfies the invariant
- Benefits:
  - **Safety**: The other parts of the program can't cause bugs in the set implementation
  - **Modularity**: It is possible to change the implementation without changing the rest of the program

# Summary: Abstract Types

- Different programming languages have different ways of letting you define abstract types
- At a minimum, this means providing:
  - A way to specify (write down) an interface
  - A means of hiding implementation details (*encapsulation*)
- In OCaml:
  - Interfaces are specified using a *signature* or *interface*
  - Encapsulation is achieved because the interface can *omit* information
    - type definitions
    - names and types of auxiliary functions
  - Clients *cannot* mention values or types not named in the interface

# Dealing with Partiality\*

\*A function is said to be *partial* if it is not defined for all inputs.

# Which of these is a function that calculates the maximum value in a (generic) list?

Which of these is a function that calculates the maximum value in a (generic) list:

1. 

```
let rec list_max (l:'a list) : 'a =  
  begin match l with  
  | [] -> []  
  | h :: t -> max h (list_max t)  
  end
```

2. 

```
let rec list_max (l:'a list) : 'a =  
  fold max 0 l
```

3. 

```
let rec list_max (l:'a list) : 'a =  
  begin match l with  
  | h :: t -> max h (list_max t)  
  end
```

4. None of the above

1  
2  
3  
4

Which of these is a function that calculates the maximum value in a (generic) list:

1. 

```
let rec list_max (l:'a list) : 'a =  
  begin match l with  
  | [] -> []  
  | h :: t -> max h (list_max t)  
  end
```

2. 

```
let rec list_max (l:'a list) : 'a =  
  fold max 0 l
```

3. 

```
let rec list_max (l:'a list) : 'a =  
  begin match l with  
  | h :: t -> max h (list_max t)  
  end
```

4. None of the above

Answer: 4

# Quiz answer

- `list_max` isn't defined for the empty list!

```
let rec list_max (l:'a list) : 'a =  
  begin match l with  
    | [] -> failwith "empty list"  
    | [h] -> h  
    | h::t -> max h (list_max t)  
  end
```

## Client of list\_max

```
(* string_of_max calls list_max *)  
let string_of_max (x:int list) : string =  
  string_of_int (list_max x)
```

- Oops! `string_of_max` will fail if given `[]`
- Not so easy to debug if `string_of_max` is written by one person and `list_max` is written by another.

- Interface of `list_max` is not very informative  
`val list_max : int list -> int`

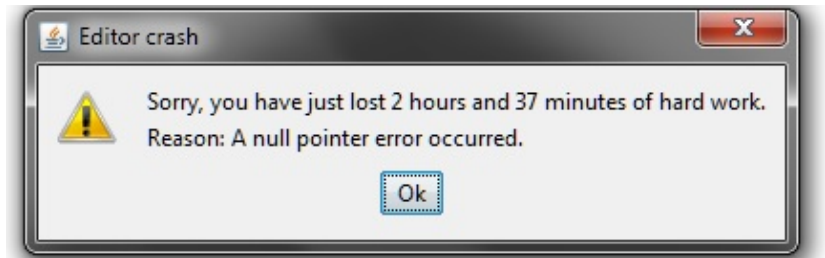
# Solutions to Partiality: Option 1

- Abort the program:
  - `failwith "an error message"`
    - Whenever it is called, `failwith` halts the program and reports the error message it is given.
- This solution is appropriate whenever you *know* that a certain case is impossible
  - The compiler isn't smart enough to figure out that the case is impossible...
  - Often happens when there is an invariant on a data structure
  - `failwith` is also useful to "stub out" unimplemented parts of your program.
- Languages (e.g. OCaml, Java) support *exception handling facilities* to let programs recover from such failures.
  - We'll talk about these when we get to Java



# Solutions to Partiality: Option 2

- Return a *default or error value*
  - e.g. define `list_max []` to be `-1`
  - Error codes used often in C programs
  - `null` used often in Java
- But...
  - What if `-1` (or whatever default you choose) really *is* the maximum value?
  - Can lead to many bugs if the default isn't handled properly by the callers.
  - *IMPOSSIBLE* to implement generically!
    - No way to generically create a sensible default value for every possible type
  - Sir Tony Hoare, Turing Award winner and inventor of `null` calls it his "*billion dollar mistake*"!
- *Defaults should be avoided if possible*



# Optional values

Solutions to Partiality: Option 3

# Option Types

- Define a generic datatype of *optional values*:

```
type 'a option =  
  | None  
  | Some of 'a
```

- A “partial” function returns an option

```
let list_max (l:list) : int option = ...
```

- Contrast this with “null”, a “legal” return value of any type
  - caller can accidentally forget to check whether null was used; results in NullPointerExceptions or crashes
- Modern language designs (e.g. Apple's Swift, Mozilla's Rust) distinguish between the type String (definitely not null) and String? (optional string)

## Example: list\_max

- A function that returns the maximum value of a list as an option (None if the list is empty)

```
let list_max (l:'a list) : 'a option =  
  begin match l with  
    | [] -> None  
    | x::tl -> Some (fold max x tl)  
  end
```

## Revised client of list\_max

```
(* string_of_max calls list_max *)  
let string_of_max (l:int list) : string =  
  begin match (list_max l) with  
  | None -> "no maximum"  
  | Some m -> string_of_int m  
  end
```

- string\_of\_max will never fail
- The type of list\_max makes it explicit that a *client* must check for partiality.

```
val list_max : int list -> int option
```

# What is the type of this function?

```
let head (x: _____) : _____ =  
  begin match x with  
  | [] -> None  
  | h :: t -> Some h  
  end
```

'a list -> 'a

'a list -> 'a list

'a list -> 'b option

'a list -> 'a option

None of the above

# Which of these is a function that calculates the maximum value in a (generic) list?

Which of these is a function that calculates the maximum value in a (generic) list:

1. 

```
let rec list_max (l:'a list) : 'a =  
begin match l with  
| [] -> []  
| h :: t -> max h (list_max t)  
end
```

2. 

```
let rec list_max (l:'a list) : 'a =  
fold max 0 l
```

3. 

```
let rec list_max (l:'a list) : 'a =  
begin match l with  
| h :: t -> max h (list_max t)  
end
```

4. None of the above

1  
2  
3  
4

What is the type of this function?

```
let head (x: _____) : _____ =  
  begin match x with  
  | [] -> None  
  | h :: t -> Some h  
  end
```

1. 'a list -> 'a
2. 'a list -> 'a list
3. 'a list -> 'b option
4. 'a list -> 'a option
5. None of the above

Answer: 4



# What is the value of this expression?

```
let head (x: 'a list) : 'a option =  
  begin match x with  
  | [] -> None  
  | h :: t -> Some h  
  end in
```

```
[ head [1]; head [] ]
```

[1; 0]

1

[Some 1; None]

[None; None]

None of the above

What is the value of this expression?

```
let head (x: 'a list) : 'a option =  
  begin match x with  
  | [] -> None  
  | h :: t -> Some h  
  end in
```

```
[ head [1]; head [] ]
```

1. [ 1 ; 0 ]
2. 1
3. [Some 1; None]
4. [None; None]
5. None of the above

Answer: 3

# Revising the MAP interface

```
module type MAP = sig

  type ('k, 'v) map

  val empty      : ('k, 'v) map
  val add        : 'k -> 'v -> ('k, 'v) map -> ('k, 'v) map
  val remove    : 'k          -> ('k, 'v) map -> ('k, 'v) map
  val mem       : 'k -> ('k, 'v) map -> bool
  val get       : 'k -> ('k, 'v) map -> 'v option
  val entries   : ('k, 'v) map -> ('k * 'v) list
  val equals    : ('k, 'v) map -> ('k, 'v) map -> bool

end
```

get returns an optional 'v.  
Now its type isn't a lie!

# Commands, Sequencing and Unit

What is the type of `print_string`?

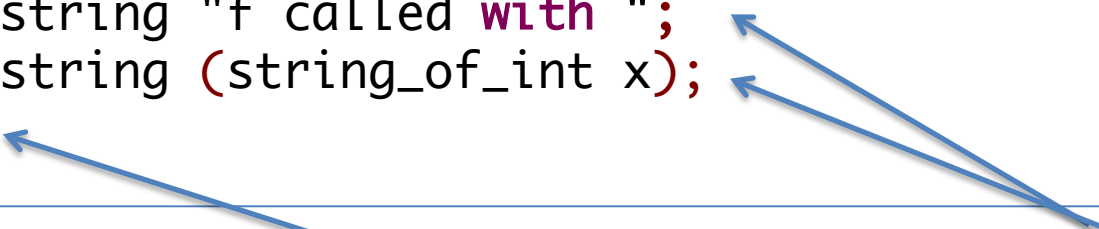


# Sequencing Commands and Expressions

We can *sequence* commands inside expressions using ‘;’

- unlike in C, Java, etc., ‘;’ doesn’t terminate a statement it *separates* a command from an expression

```
let f (x:int) : int =  
  print_string "f called with ";  
  print_string (string_of_int x);  
  x + x
```



do not use ‘;’ here!

note the use of ‘;’ here

The distinction between commands & expressions is artificial.

- `print_string` is a function of type: `string -> unit`
- Commands are actually just expressions of type: `unit`

# unit: the trivial type

- Similar to "void" in Java or C
- For functions that don't take any arguments

```
let f () : int = 3
let y : int = f ()
```

```
val f : unit -> int
val y : int
```

- Also for functions that don't return anything, such as testing and printing functions a.k.a *commands*:

```
(* run_test : string -> (unit -> bool) -> unit *)
;; run_test "TestName" test

(* print_string : string -> unit *)
;; print_string "Hello, world!"
```

# unit: the boring type

- *Actually, () is a value just like any other value (a 0-ary tuple)*
- For functions that don't take any **interesting** arguments

```
let f () : int = 3
let y : int = f ()
```

```
val f : unit -> int
val y : int
```

- Also for functions that don't return anything **interesting**, such as testing and printing functions a.k.a *commands*:

```
(* run_test : string -> (unit -> bool) -> unit *)
;; run_test "TestName" test

(* print_string : string -> unit *)
;; print_string "Hello, world!"
```

# unit: the first-class type

- Can define values of type unit

```
let x : unit = ()
```

```
val x : unit
```

- Can pattern match unit (even in function definitions)

```
let z = begin match x with  
  | () -> 4  
end
```

```
fun () -> 3
```

- Is the result of an implicit else branch:

```
;; if z <> 4 then  
  failwith "oops"
```

=

```
;; if z <> 4 then  
  failwith "oops"  
else ()
```



# Sequencing Commands and Expressions

- Expressions of type unit are useful because of their *side effects* – they "do" stuff
  - e.g. printing, changing the value of mutable state

```
let f (x:int) : int =  
  print_string "f called with ";  
  print_string (string_of_int x);  
  x + x
```

do not use ';' here!

note the use of ';' here

- We can think of ';' as an infix function of type:  
unit -> 'a -> 'a

# What is the type of f in the following program?

```
let f (x:int) =  
  print_int (x + x)
```

unit -> int

unit -> unit

int -> unit

int -> int

f is ill typed

What is the type of `f` in the following program:

```
let f (x:int) =  
    print_int (x + x)
```

1. `unit -> int`
2. `unit -> unit`
3. `int -> unit`
4. `int -> int`
5. `f` is ill typed

Answer: 3

# What is the type of f in the following program?

```
let f (x:int) =  
  (print_int x);  
  (x + x)
```

unit -> int

unit -> unit

int -> unit

int -> int

f is ill typed

What is the type of `f` in the following program:

```
let f (x:int) =  
  (print_int x);  
  (x + x)
```

1. `unit -> int`
2. `unit -> unit`
3. `int -> unit`
4. `int -> int`
5. `f` is ill typed

Answer: 4