

Readings

- [Lecture Notes Chapter 15: Huffman Coding](#)

Review: Huffman Coding

The motivation behind Huffman Coding is to encode and decode characters as bits, minimizing the average bits per letter (ABL). Furthermore, we seek a prefix-free code, where no encoding is a prefix of another — implying that a bit sequence can be parsed and decoded without any ambiguity.

The Huffman algorithm is a greedy algorithm that does this. Given a set of characters and their frequencies, the algorithm outputs an encoding by repeatedly merging the 2 nodes with the smallest frequency values until only one node remains. This one node is the root of the Huffman tree, whose leaves are characters and each root-to-leaf path is an encoding of that character. Furthermore, this tree is a full binary tree, where each internal node has exactly 2 children. Therefore, the Huffman algorithm produces an optimal and prefix-free encoding that minimizes the ABL.

The running time of the Huffman algorithm is $O(n \log n)$ if we utilize a min-heap to find the 2 nodes with minimum frequency in each step, as seen in the [pseudocode](#). This is because at each step, we perform a constant number of EXTRACT-MIN and INSERT operations, which take $O(\log n)$ time, and we repeat this for $O(n)$ iterations.

Problems

Problem 1

Construct an optimal Huffman coding for the following alphabet and frequency table S :

Character:	A	B	C	D	E
Frequency:	0.4	0.3	0.15	0.1	0.05

What is the ABL, or average bits per letter, for this encoding?

Problem 2

You have an alphabet with $n > 2$ letters and frequencies. You perform Huffman encoding on this alphabet, and notice that the character with the largest frequency is encoded by just a 0. In this alphabet, symbol i occurs with probability p_i ; $p_1 \geq p_2 \geq p_3 \geq \dots \geq p_n$.

Given this alphabet and encoding, does there exist an assignment of probabilities to p_1 through p_n such that $p_1 < \frac{1}{3}$? Justify your answer.