

Lecture 4

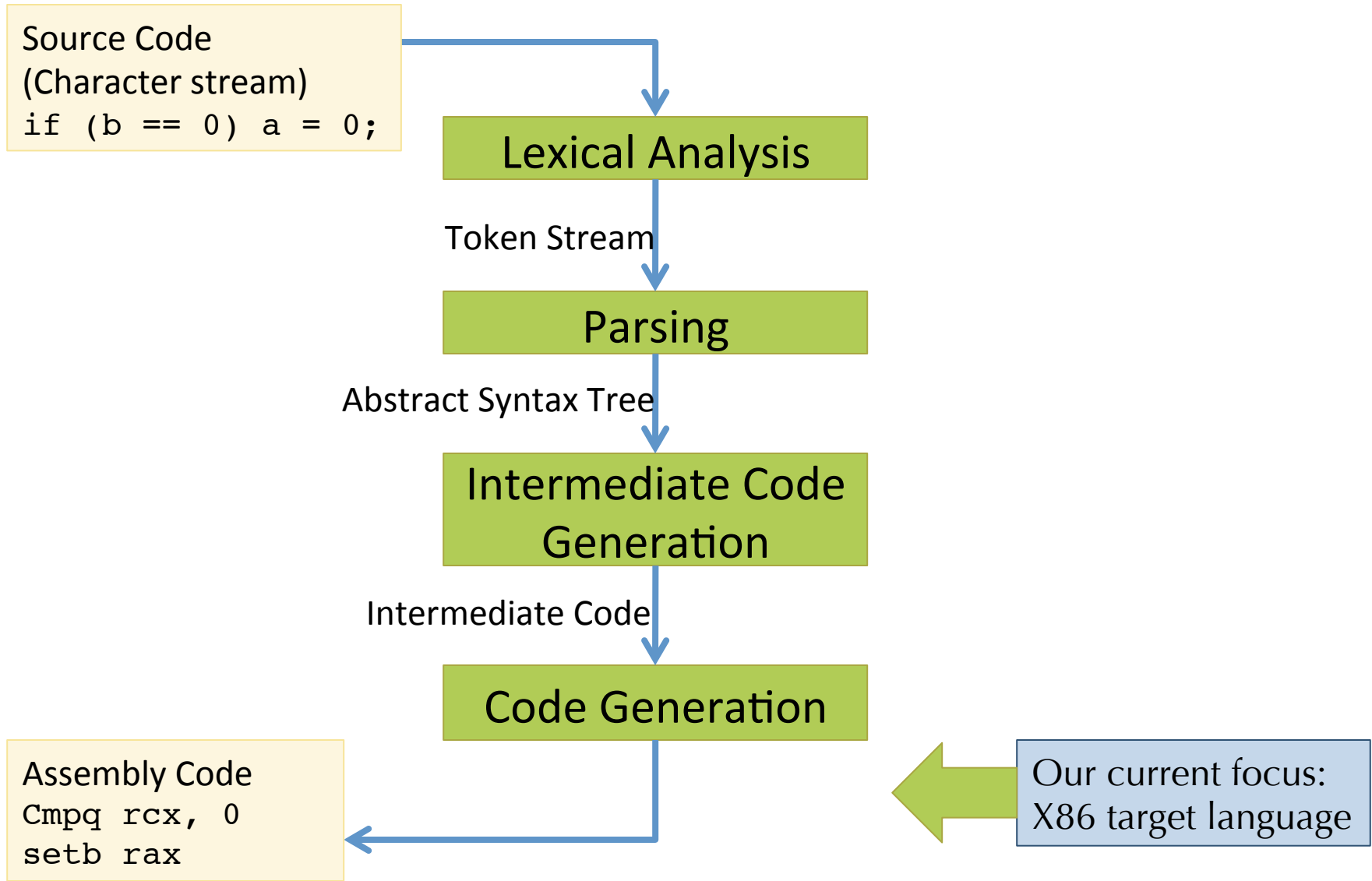
CIS 341: COMPILERS

Announcements

- HW2: X86lite
 - Available on the course web pages.
 - Due: Monday, February 2nd at 11:59:59pm
 - Pair-programming:
 - There's a pair-search survey on Piazza
 - Register the group on the submission page
 - Submission by any group member counts for the group

- Registration:
 - If you were on the wait list, you should have been contacted
 - If you are *not* registered, please see me after class

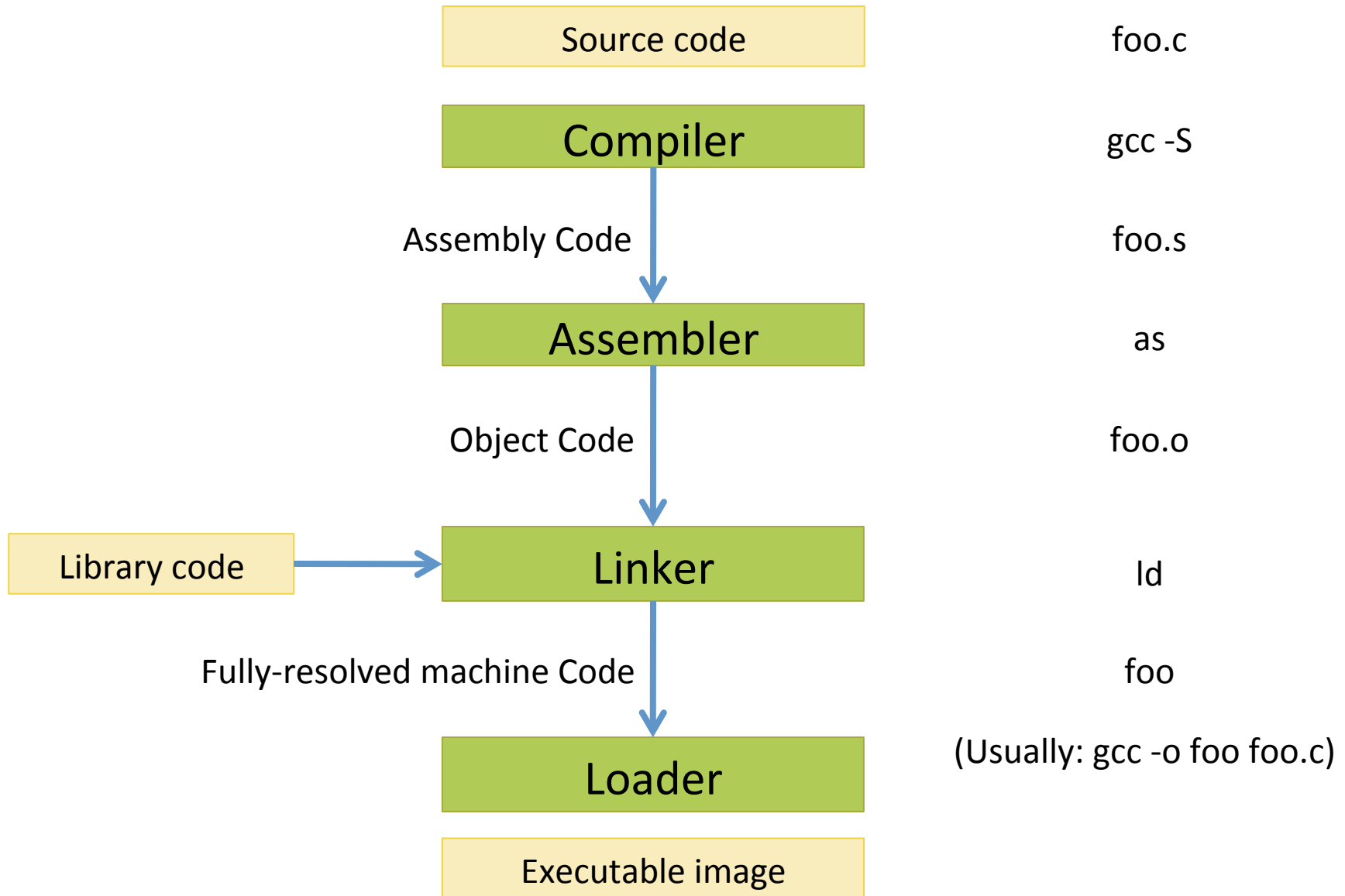
(Simplified) Compiler Structure



See www.cis.upenn.edu/~cis341/15sp/hw/hw2/x86lite.shtml

X86LITE

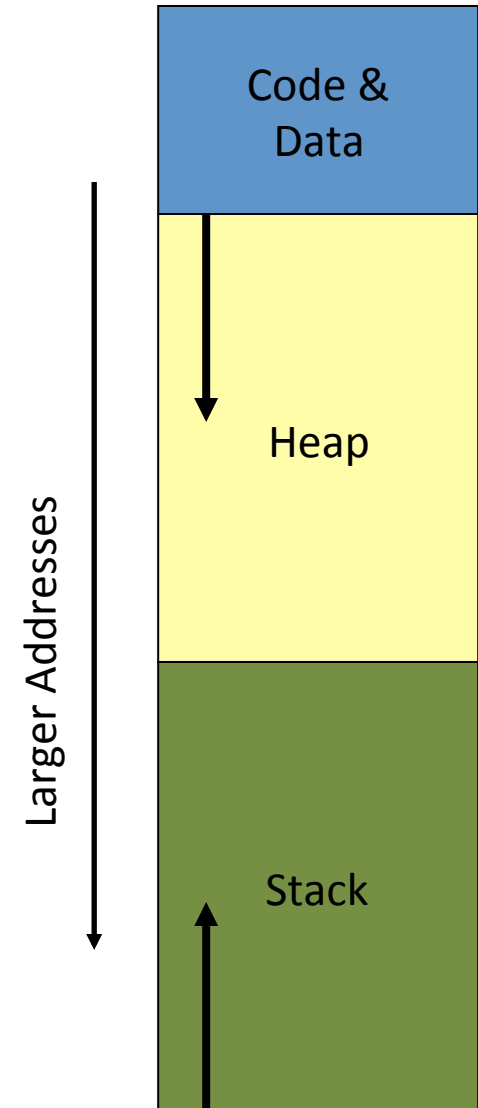
Compilation & Execution



PROGRAMMING IN X86LITE

3 parts of the C memory model

- The code & data (or "text") segment
 - contains compiled code, constant strings, etc.
- The Heap
 - Stores dynamically allocated objects
 - Allocated via "malloc"
 - Deallocated via "free"
 - C runtime system
- The Stack
 - Stores local variables
 - Stores the return address of a function
- In practice, most languages use this model.



Local/Temporary Variable Storage

- Need space to store:
 - Global variables
 - Values passed as arguments to procedures
 - Local variables (either defined in the source program or introduced by the compiler)
- Processors provide two options
 - Registers: fast, small size (32 or 64 bits), limited number
 - Memory: slow, very large amount of space (4+ GB)
- In practice on X86:
 - Registers are limited (and have restrictions)
 - Divide memory into regions including the *stack* and the *heap*

Calling Conventions

- Specify the locations (e.g. register or stack) of arguments passed to a function
- Designate registers either:
 - Caller Save – e.g. freely usable by the called code
 - Callee Save – e.g. must be restored by the called code
- Protocol for deallocating stack-allocated arguments
 - Caller cleans up
 - Callee cleans up (makes variable arguments harder)

32-bit cdecl calling conventions

- “Standard” on X86 for many C-based operating systems (i.e. almost all)
 - Still some wrinkles about return values (e.g. some compilers use **EAX** and **EDX** to return small values)
 - This is evolving due to 64 bit (which allows for packing multiple values in one register)
- Arguments are passed on the stack in right-to-left order
- Return value is passed in **EAX**
- Registers **EAX**, **ECX**, **EDX** are caller save
- Other registers are callee save
 - Ignoring these conventions will cause havoc (bus errors or seg faults)

- Many other variants: fastcall, syscall, thiscall

x86-64 calling conventions

- Microsoft x64
 - Used by Visual C++ and Windows (but supported by gcc, intel C++, etc.)
 - 4 register arguments
 - 4-quad “shadow space”
- System V AMD64 ABI
 - Used by linux, bsd, Mac OSX
 - First six integer/pointer arguments are passed in registers:
 - rdi, rsi, rdx, rcx, r8, r9
 - Arguments seven and up, passed on the stack
 - Stack aligned on 16-byte boundaries
 - Callee save registers: rbp, rbx, r12—r15
 - Caller save register: everything else
 - Caller cleans up stack arguments

Call Stacks: Caller's protocol

- Function call:

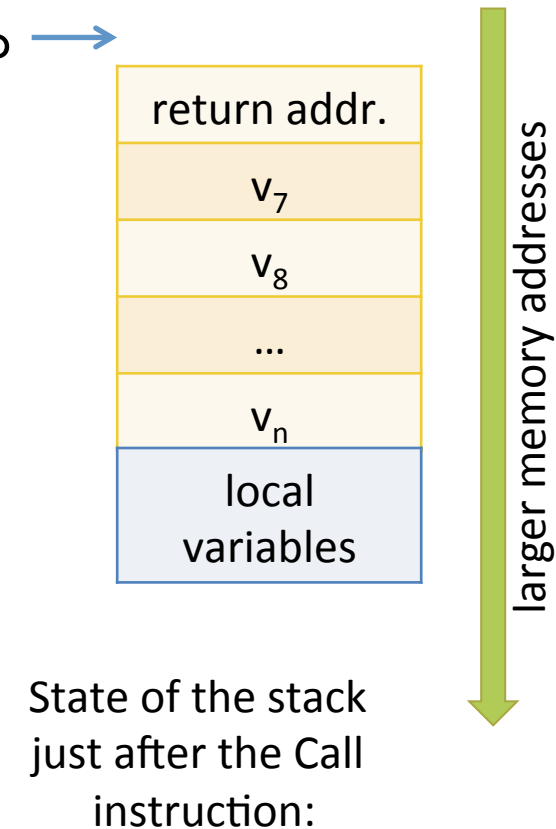
$f(e_1, e_2, \dots, e_n);$

- Save caller-save registers:
 - all but `rbp`, `rbx`, `r12-r15`
- Evaluate e_1 to v_1 , e_2 to v_2 , ..., e_n to v_n
- Move $v_1 .. v_6$ into registers as on previous slide.
- Push v_7 to v_n onto the top of the stack.
- Use `call` to jump to the code for f
 - pushing the return address onto the stack.

- Invariant: returned value passed in `rax`

- After call:

- clean up the pushed arguments by popping the stack.
- Restore caller-saved registers

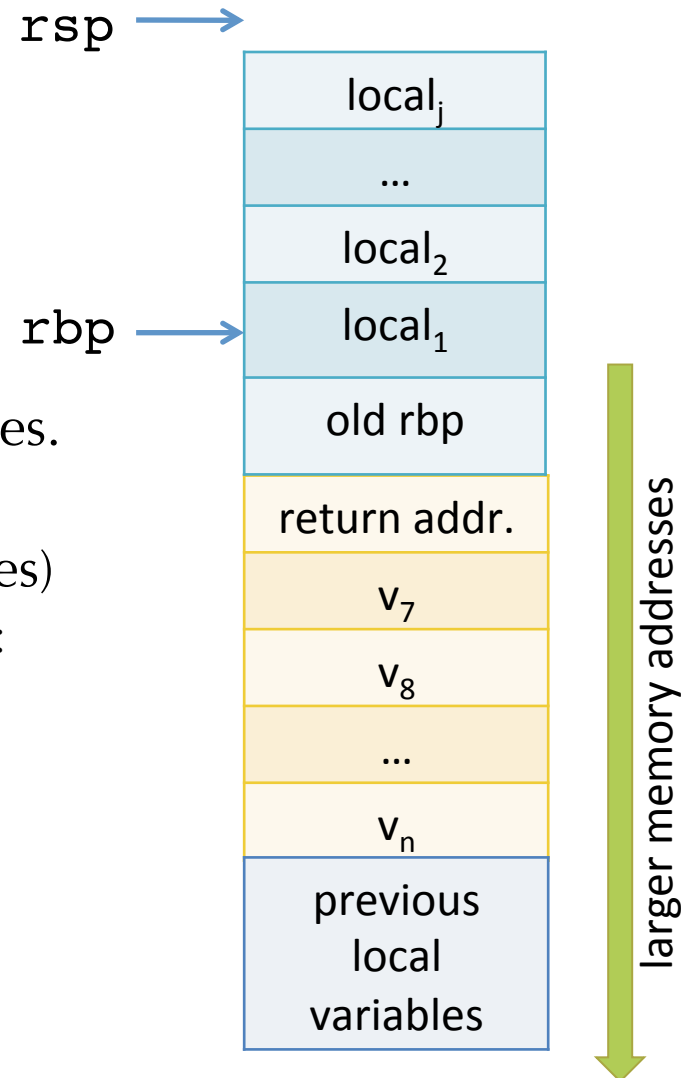


Call Stacks: Callee's protocol

- On entry:
 1. Save old frame pointer
 - rbp is callee save
 2. Create new frame pointer
 - `movq rsp, rbp`
 3. Allocate stack space for local variables.

- Invariants: (assuming quad-size values)
 - Function argument $n > 6$ is located at:
 $rbp + (n-5) * 8$
 - Local variable $local_j$ is located at:
 $rbp - (j - 1) * 8$

- On exit:
 1. Pop local storage
 2. Restore rbp



State of the stack
after Step 3 of entry.