Introductions, C refresher

Computer Operating Systems, Fall 2023

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How are you?

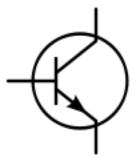
Lecture Outline

Introduction & Logistics

- Course Overview
- Assignments & Exams
- Policies
- C Refresher
 - Memory Layout
 - Demo (make, man pages)
 - Malloc, free, pointers
 - stdin, stdout

Instructor: Travis McGaha

- UPenn CIS faculty member since August 2021
 - First Semester with CIS 3800
 - CIS 2400 in 21fa & 22fa
 - CIT 5950 in 22sp & 23sp
- More on my personal website: https://www.cis.upenn.edu/~tqmcgaha/
- Schedule meeting w/ me
- Unofficial office hours right after class
- Official office hours TBD





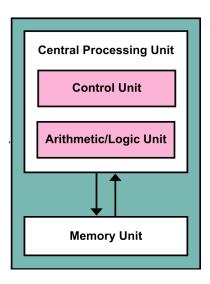




Adder

Mux/Demux

Latch/Flip-Flop



Process

Operating System

Computer

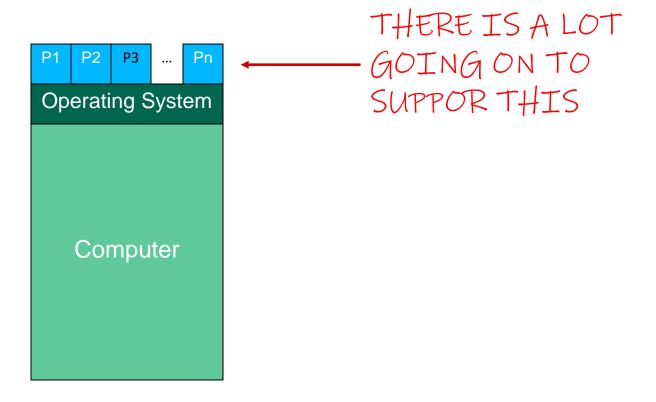


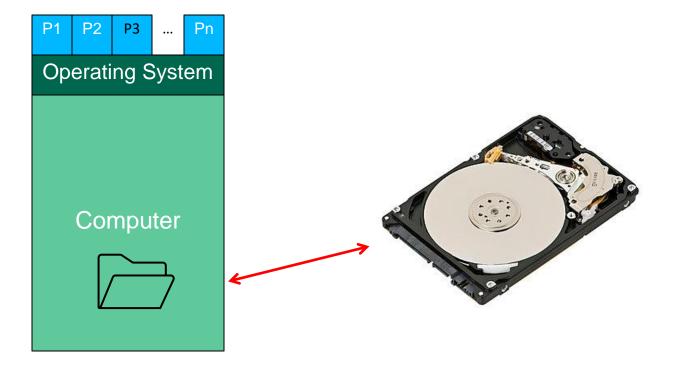
Process

Operating System

Computer

OS does A LOT more than just printing, reading input, video display, and timer





Prerequisites

- Course Prerequisites:
 - CIS 2400
 - Teamwork & Willingness/happy to spend substantial time coding
- What you should be familiar with already:
 - C programming
 - C Memory Model
 - Computer Architecture Model
 - Basic UNIX command line skills
- HW0 is tuned so that it will help refresh you on these.
 - But it still covers new content!
 - Even if you think you know C, get started sooner rather than later. 17

CIS 3800 Learning Objectives

- To leave the course with a better understanding of:
 - How a computer runs/manages multiple programs
 - How the previous point may affect the code we write
 - How to read documentation
 - Experience writing a massive programming project FROM SCRATCH with others.
 - More comfortable writing C code
- Topics list/schedule can be found on course website
 - Note: These topics may be tweaked

Disclaimer

- This is a digest, <u>READ THE SYLLABUS</u>
 - https://www.seas.upenn.edu/~cis3800/current/documents/syllab
 us

Course Components: Textbook

- Textbook (0)
 - Textbooks recommended in pasts
 - A.S. Tanenbaum. Modern Operating Systems (4th Edition onwards).
 Prentice-Hall.
 - W. Richard Stevens and Stephen A. Rago. Advanced Programming in the UNIX Environment (2/e or 3/e). Addison-Wesley Professional.
 - Systems for all: https://diveintosystems.org/book/
 - Free online textbook, pretty well written
 - Linux Man pages:
 - https://linux.die.net/man/
 - https://www.man7.org/linux/man-pages/
 - The man command in the terminal
 - DEMO:
 - name a C function
 - tcsetpgrp

21

Course Components: Part 1

- Lectures (~26)
 - Introduces concepts, slides & recordings available on canvas
 - In lecture polling. Polls are not counted towards credit
- Pre-recorded videos (many)
 - Entirely optional
 - Goes over lecture material or demonstrates something for projects
- Check-ins "Quizzes" (~12)
 - Unlimited attempt low-stake quizzes on canvas to make sure you are caught up with material
 - Lowest two are dropped
- Midterms (2)
 - Two in-person exams, two pages of notes allowed
 - Details TBD

Programming Facilities

- Docker
 - Same environment as the autograder
 - Instructions for setup to be posted soon
- Speclab cluster, as a fallback incase Docker does not work
 - Instructions on course website
 - To see status: https://www.seas.upenn.edu/checklab/?lab=speclab
- DO NOT use Eniac machines to develop projects for this class!

Project 0 & 1

Project 0

- Unix "Shell" command interpreter (e.g. sh, bash, etc)
- Excellent way to learn about how system calls are supported and used.
- Done individually
- Code review
- Will be posted soon!!

Project 1

- Unix "Shell" the real deal
- Redirection, pipelines, background/foreground processing, job control
- Groups of two.

PennOS

- Best way to learn about an operating systems is to build one.
- Build all the main features of an OS (in emulation)
- Groups of 4.
- By the end of the project, you will:
 - Learn about how different subsystems in Unix interact with each other
 - Learn about priority scheduling, file systems, user shell interactions
 - Become a really good and confident systems programmer

PennOS

- There is a paper on this: http://netdb.cis.upenn.edu/papers/pennos.pdf at an ACM OS journal.
- Group evaluation done by the end of semester. Team member with lower than 15% contribution to the group will get grade downgrade.

HW Policies

 Students who did not contribute to group projects will get F grade regardless of overall score.

Late Policy

- You are given 5 late tokens.
- Tokens are counted per student and can only be used on some assignments.
- Two tokens used at max per assignment
- Each token grants 48 hours of extra time
- If there are extenuating circumstances, please let me know

Collaboration Policy Violation

- You will be caught:
 - Careful grading of all written homeworks by teaching staff
 - Measure of Software Similarity (MOSS): http://theory.stanford.edu/~aiken/moss/
 - Successfully used in several classes at Penn
- Zero in assignment, zero for class participation (3%). F grade if caught twice.
 - First-time offenders will be reported to Office of Student Conduct with no exceptions. Possible suspension from school
 - Your friend from last semester who gave the code will have their grade retrospectively downgraded.

Collaboration Policy Violation

- Generative Al
 - I am skeptical of its usefulness for your learning and for your success in the course
 - Some articles on the topic:
 - https://www.aisnakeoil.com/p/chatgpt-is-a-bullshit-generator-but
 - https://www.aisnakeoil.com/p/gpt-4-and-professional-benchmarks
 - Not banned, but not recommended. Use your best judgement.
- You will not help your overall grade and happiness:
 - Quizzed individually during project demo, exams on project in finals
 - If you can't explain your code in OH, we can turn you away.
 - This is different than being confused on a bug or with C, this is ok
 - Personal lifelong satisfaction from completing PennOS

Course Grading

- Breakdown:
 - Project 0 penn-shredder: (8%)
 - Project 1 penn-shell: (18%)
 - Project 2 PennOS: (37%)
 - Exams (34%)
 - 17% each
 - Check-in Quizzes(3%)
- Final Grade Calculations:
 - I would LOVE to give everyone an A+ if it is earned
 - Final grade cut-offs will be decided privately at the end of the Semester. What is used in previous semester is in the syllabus

Course Infrastructure

- Course Website: www.seas.upenn.edu/~cis3800/23fa/
 - Materials, Schedule, Syllabus ...
- Docker or Speclab
 - Coding environment for hw's
- Gradescope
 - Used for HW Submissions
- Poll Everywhere
 - Used for lecture polls
- Ed Discussion
 - Course discussion board

Getting Help

Ed

- Announcements will be made through here
- Ask and answer questions
- Sign up if you haven't already!

Office Hours:

- Can be found on calendar on front page of course website
- Starts next week for all TAs

❖ 1-on-1's:

- Can schedule 1-on-1's with Travis
- Should attend OH and use Ed when possible, but this is an option for when OH and Ed can't meet your needs

We Care

- I am still figuring things out, but we do care about you and your experience with the course
 - Please reach out to course staff if something comes up and you need help

PLEASE DO NOT CHEAT OR VIOLATE ACADEMIC INTEGRITY

- We know that things can be tough, but please reach out if you feel tempted. We want to help
- Read more on academic integrity in the syllabus



Poll Everywhere

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Any questions, comments or concerns so far?

Lecture Outline

- Introduction & Logistics
 - Course Overview
 - Assignments & Exams
 - Policies
- C refresher
 - Pointers
 - Arrays

I Will go through parts of this relatively fast.

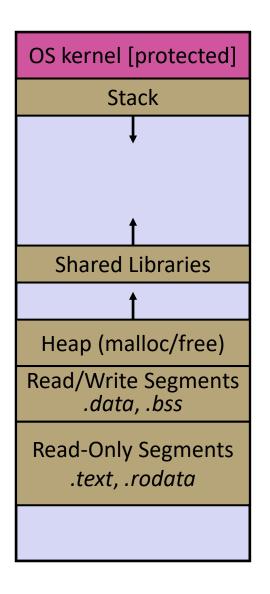
Review this on your own





Aside: Memory

- Where all data, code, etc are stored for a program
- Broken up into several segments:
 - The stack
 - The heap
 - The kernel
 - Etc.
- Each "unit" of memory has an address



Pointers

POINTERS ARE EXTREMELY IMPORTANT IN C

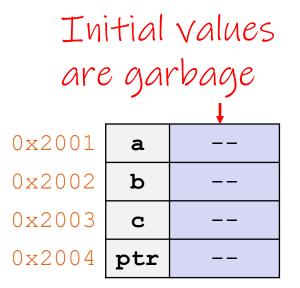
- Variables that store addresses
 - It stores the address to somewhere in memory
 - Must specify a type so the data at that address can be interpreted
- Generic definition: type* name; or type *name;

 Example: int *ptr;
 - Declares a variable that can contain an address
 - Trying to access that data at that address will treat the data there as an int

Pointer Operators

- Dereference a pointer using the unary * operator
 - Access the memory referred to by a pointer
 - Can be used to read or write the memory at the address
 - Example: int *ptr = ...; // Assume initialized int a = *ptr; // read the value *ptr = a + 2; // write the value

- Get the address of a variable with &
 - &foo gets the address of foo in memory
 - Example: int a = 595;
 int *ptr = &a;
 *ptr = 2; // 'a' now holds 2





0x2001	a	5
0x2002	b	3
0x2003	С	
0x2004	ptr	

0x2001	a	5	
0x2002	b	3	
0x2003	С		
0x2004	ptr	0x2001	/

CIS 3800, Fall 2023



```
int main(int argc, char** argv) {
  int a, b, c;
  int* ptr;  // ptr is a pointer to an int

a = 5;
  b = 3;
  ptr = &a;

  *ptr = 7;
  c = a + b;

return 0;
}
```

0x2001	a	7	
0x2002	b	3	
0x2003	С		
0x2004	ptr	0x2001	/



```
int main(int argc, char** argv) {
  int a, b, c;
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a = 5;
  b = 3;
  ptr = &a;

*ptr = 7;
  c = a + b;

return 0;
}
```

0x2001	a	7	
0x2002	b	3	
0x2003	C	10	/
0x2004	ptr	0x2001	/



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What does this code print?

```
#include <stdio.h>
#include <stdlib.h>
void modify_int(int x) {
  x = 5;
int main() {
  int num = 3;
  modify int(num);
  printf("%d\n", num);
  return EXIT SUCCESS;
```

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What does this code print?

How could we fix it?
 E.g. make modify point actually modify a point

```
#include <stdio.h>
#include <stdlib.h>
typedef struct point st {
  int x, y;
} Point;
void modify point(Point p) {
  p.x = 3800;
  p.y = 4710
int main() {
  Point p = \{1100, 2400\};
  modify point(p);
  printf("%d, %d\n", p.x, p.y);
  return EXIT SUCCESS;
```

Demo: pass_by.c

- Everything in C is pass-by value (e.g. a copy is passed to the function)
- HOWEVER, we can pass a copy of a pointer (e.g. a reference to something) to mimic pass-by-reference.
- Demo pass_by.c
 - Note: most lecture code will be available on the course website

 Pointers can be used to "return" more than one value from a function

```
int solve quadratic (double a, double b, double c,
                    double* soln1, double* soln2) {
  double d = b*b - 4 * a * c;
  if (d >= 0) {
    *soln1 = (-b + sqrt(d)) / (2*a);
    *soln2 = (-b - sqrt(d)) / (2*a);
    return 1;
  } else {
   return 0;
int main(int argc, char** argv) {
  double soln1, soln2; // populated by function call
  solve quadratic(2.0, 4.0, 0.0, &soln1, &soln2);
  // . . .
```

 Pointers can be used to "return" more than one value from a function

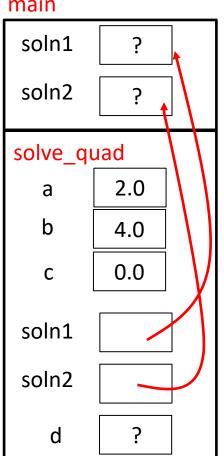
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    return 1;
  } else {
   return 0;
int main(int argc, char** argv) {
 double soln1, soln2; // populated by function call
 solve quadratic(2.0, 4.0, 0.0, &soln1, &soln2);
```

main

```
soln1
soln2
```

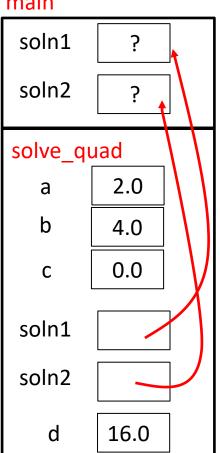
 Pointers can be used to "return" more than one value from a function main

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  // ...
```



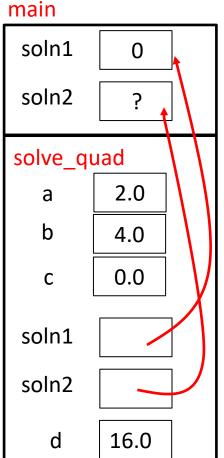
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  solve quadratic(2.0, 4.0, 0.0, &soln1, &soln2);
  // ...
```



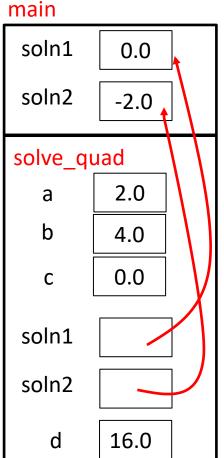
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  // ...
```



Pointers can be used to "return" more than one value from a function

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    *soln2 = (-b - sqrt(d)) / (2*a);
 \rightarrow return 1:
 } else {
   return 0;
int main(int argc, char** argv) {
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  // ...
```



 Pointers can be used to "return" more than one value from a function

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    return 1;
  } else {
   return 0;
int main(int argc, char** argv) {
  double soln1, soln2; // populated by function call
  solve quadratic(2.0, 4.0, 0.0, &soln1, &soln2);
```

main

```
soln1
          0.0
soln2
          -2.0
```

Arrays

- Definition: [type name[size]
 - Allocates size*sizeof (type) bytes of contiguous memory
 - Normal usage is a compile-time constant for size (e.g. int scores[175];)
 - Initially, array values are "garbage"

- Size of an array
 - Not stored anywhere array does not know its own size!
 - The programmer will have to store the length in another variable or hard-code it in

Using Arrays

Optional when initializing

- hitialization: type name[size] = {val0,...,valN};
 - { } initialization can *only* be used at time of definition
 - If no size supplied, infers from length of array initializer
- Array name used as identifier for "collection of data"
 - name [index] specifies an element of the array and can be used as an assignment target or as a value in an expression
 - Array name (by itself) produces the address of the start of the array
 - Cannot be assigned to / changed

```
int primes[6] = {2, 3, 5, 6, 11, 13};
primes[3] = 7;
primes[100] = 0; // memory smash!
```

No IndexOutOfBounds Hope for seafault

Multi-dimensional Arrays

Generic 2D format:

```
type name[rows][cols];
```

- Still allocates a single, contiguous chunk of memory
- C is row-major
- Can access elements with multiple indices

```
A[0][1] = 7;my int = A[1][2];
```

The entries in this array are stored in memory in row major order as follows:

```
•A[0][0], A[0][1], A[0][2], A[1][0], A[1][1], A[1][2]
```

2-D arrays normally only useful if size known in advance.
 Otherwise use dynamically-allocated data and pointers (later)

Arrays as Parameters

- It's tricky to use arrays as parameters
 - What happens when you use an array name as an argument?
 - Arrays do not know their own size

Passes in address of start of array

```
int sumAll(int a[]) {
  int i, sum = 0;
  for (i = 0; i < ...???
}</pre>
```

```
int sumAll(int* a) {
  int i, sum = 0;
  for (i = 0; i < ...???
}</pre>
```

Equivalent

- Note: Array syntax works on pointers
 - E.g. [ptr[3] = ...;

Solution: Pass Size as Parameter

```
int sumAll(int* a, int size) {
  int i, sum = 0;
  for (i = 0; i < size; i++) {
    sum += a[i];
  }
  return sum;
}</pre>
```

Standard idiom in C programs

Strings without Objects

- Strings are central to C, very important for I/O
- In C, we don't have Objects but we need strings
- If a string is just a sequence of characters, we can have use array of characters as a string

Example:

```
char str_arr[] = "Hello World!";
char *str_ptr = "Hello World!";
```

Null Termination

DO NOT FORGET THIS. THIS IS THE CAUSE OF MANY BUGS

- Arrays don't have a length, but we mark the end of a string with the null terminator character.
 - The null terminator has value 0x00 or '\0'
 - Well formed strings <u>MUST</u> be null terminated
- * Example: char str[] = "Hello";
 - Takes up 6 characters, 5 for "Hello" and 1 for the null terminator

address	0x2000	0x2001	0x2002	0x2003	0x2004	0x2005
value	'H'	'e'	'1'	'1'	'0'	'\0'

Demo: get_input.c

- Lets code together a small program that:
 - Reads at max 100 characters from stdin (user input)
 - Truncates the input to only the first word
 - Prints that word out
 - Not allowed to use scanf, FILE*, printf, etc

Poll Everywhere

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- There is something wrong with this function
- What is it? How do we fix this function w/o changing the function signature

```
#define MAX INPUT SIZE 100
char* read stdin() {
  char str[MAX INPUT SIZE];
  ssize t res = read(STDIN FILENO, str, MAX INPUT SIZE);
  // error checking
  <u>if</u> (res <= 0) {
    return NULL;
  return str;
```

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The Stack

main	
char* result	
	i
	! !
	! !
	l

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#define MAX INPUT_SIZE 100
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```



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The Stack

```
main
char* result

read_stdin
str ['H', 'i', '\0']
```

```
#define MAX INPUT_SIZE 100
char* read stdin() {
  char str[MAX INPUT SIZE];
  ssize t res = read(STDIN FILENO,
                      str, MAX INPUT SIZE);
  // error checking
  if (res <= 0) {
    return NULL;
  return str;
```

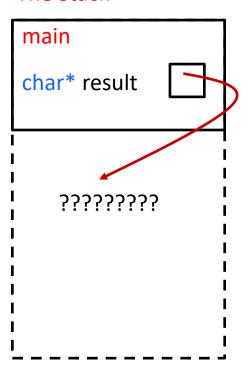
Poll Everywhere

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The Stack



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                      str, MAX INPUT SIZE);
  // error checking
  if (res <= 0) {
    return NULL;
  return str;
```

static function variables

Functions can declare a variable as static

```
#include <stdio.h> // for printf
#include <stdlib.h> // for EXIT SUCCESS
                                     This is how some functions
int next num();
                                     (like one in projo) can
                                    "remember" things.
int main(int argc, char** argv) {
  printf("%d\n", next num()); // prints 1
  printf("%d\n", next num()); // then 2
 printf("%d\n", next num()); // then 3
  return EXIT SUCCESS;
int next num() {
  // marking this variable as static means that
  // the value is preserved between calls to the function
  // this allows the function to "remember" things
  static int counter = 0;
                                      Can be thought of as a
  counter++;
                                      global variable that is
  return counter;
                                      "private" to a function
```

Memory Allocation

So far, we have seen two kinds of memory allocation:

```
int counter = 0;  // global var

int main() {
  counter++;
  printf("count = %d\n", counter);
  return 0;
}
```

- counter is statically-allocated
 - Allocated when program is loaded
 - Deallocated when program exits

```
int foo(int a) {
   int x = a + 1;  // local var
   return x;
}
int main() {
   int y = foo(10);  // local var
   printf("y = %d\n",y);
   return 0;
}
```

- a, x, y are automaticallyallocated
 - Allocated when function is called



Deallocated when function returns

Aside: NULL

- ❖ NULL is a memory location that is guaranteed to be invalid
 - In C on Linux, NULL is 0×0 and an attempt to dereference NULL causes a segmentation fault
- Useful as an indicator of an uninitialized (or currently unused) pointer or allocation error
 - It's better to cause a segfault than to allow the corruption of memory!

```
int main(int argc, char** argv) {
  int* p = NULL;
  *p = 1; // causes a segmentation fault
  return EXIT_SUCCESS;
}
```

Aside: sizeof

- * sizeof operator can be applied to a variable or a type and it evaluates to the size of that type in bytes
- Examples:
 - sizeof(int) returns the size of an integer
 - sizeof (double) returns the size of a double precision number
 - struct my_struct s;
 - sizeof(s) returns the size of the struct s
 - my_type *ptr
 - sizeof (*ptr) returns the size of the type pointed to by ptr
- Very useful for Dynamic Memory

What is Dynamic Memory Allocation?

- We want Dynamic Memory Allocation
 - Dynamic means "at run-time"
 - The compiler and the programmer don't have enough information to make a final decision on how much to allocate
 - Your program explicitly requests more memory at run time
 - The language allocates it at runtime, maybe with help of the OS
- Dynamically allocated memory persists until either:
 - A garbage collector collects it (automatic memory management)
 - Your code explicitly deallocates it (manual memory management)
- C requires you to manually manage memory
 - More control, and more headaches

Heap API

- Dynamic memory is managed in a location in memory called the "Heap"
 - The heap is managed by user-level runetime library (libc)
 - Interface functions found in <stdlib.h>
- Most used functions:
 - void *malloc(size_t size);
 - Allocates memory of specified size
 - void free(void *ptr);
 - Deallocates memory
- Note: void* is "generic pointer". It holds an address, but doesn't specify what it is pointing at.
- Note 2: size t is the integer type of sizeof()

malloc()

```
void *malloc(size_t size);
```

- malloc allocates a block of memory of the requested size
 - Returns a pointer to the first byte of that memory
 - And returns NULL if the memory allocation failed!
 - You should assume that the memory initially contains garbage
 - You'll typically use sizeof to calculate the size you need

```
// allocate a 10-float array
float* arr = malloc(10*sizeof(float));
if (arr == NULL) {
   return errcode;
}
... // do stuff with arr
```

free()

```
    Usage: free (pointer);
```

- Deallocates the memory pointed-to by the pointer
 - Pointer <u>must</u> point to the first byte of heap-allocated memory (i.e. something previously returned by malloc)
 - Freed memory becomes eligible for future allocation
 - free (NULL); does nothing.
 - The bits in the pointer are not changed by calling free
 - Defensive programming: can set pointer to NULL after freeing it

The Heap

- The Heap is a large pool of available memory to use for Dynamic allocation
- This pool of memory is kept track of with a small data structure indicating which portions have been allocated, and which portions are currently available.

* malloc:

- searches for a large enough unused block of memory
- marks the memory as allocated.
- Returns a pointer to the beginning of that memory

* free:

- Takes in a pointer to a previously allocated address
- Marks the memory as free to use.

Dynamic Memory Example

```
#include <stdlib.h>
int main() {
  char* ptr = malloc(4*sizeof(char));
  if (ptr == NULL)
    return EXIT FAILURE;
                // do stuff with ptr
  free (ptr);
```

addr	var	value
0x2001	ptr	
	• • •	1
0x4000	HEAP START	USED
0x4001		USED
0x4002		
0x4003		
0x4004		
0x4005		
0x4006		
0x4007		
0x4008		USED
0x4009		USED

Dynamic Memory Example

addr	var	value
0x2001	ptr	0x4002
• • •		
0x4000	HEAP START	USED
0x4001		USED
0x4002		USED
0x4003		USED
0x4004		USED
0x4005		USED
0x4006		
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0x4006		
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0x4008		USED
0x4009		USED



Fixed read_stdin()

```
#define MAX INPUT SIZE 100
char* read stdin() {
  char str = (char*) malloc(sizeof(char) * MAX INPUT SIZE);
  if (str == NULL) {
   return NULL:
  ssize t res = read(STDIN FILENO, str, MAX INPUT SIZE);
  // error checking
  if (res <= 0) {</pre>
   return NULL;
  return str;
```

Dynamic Memory Pitfalls

- Buffer Overflows
 - E.g. ask for 10 bytes, but write 11 bytes
 - Could overwrite information needed to manage the heap
 - Common when forgetting the null-terminator on malloc'd strings
- Not checking for NULL
 - Malloc returns NULL if out of memory
 - Should check this after every call to malloc
- Giving free () a pointer to the middle of an allocated region
 - Free won't recognize the block of memory and probably crash
- Giving free() a pointer that has already been freed
 - Will interfere with the management of the heap and likely crash
- malloc does NOT initialize memory
 - There are other functions like calloc that will zero out memory

Memory Leaks

- The most common Memory Pitfall
- What happens if we malloc something, but don't free it?
 - That block of memory cannot be reallocated, even if we don't use it anymore, until it is freed
 - If this happens enough, we run out of heap space and program may slow down and eventually crash

Garbage Collection

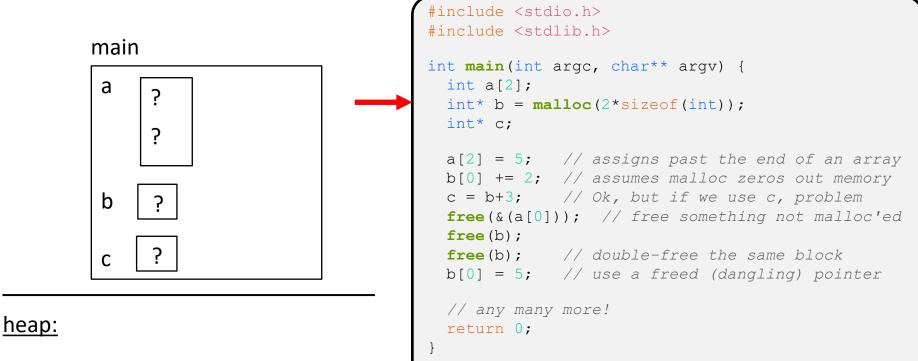
- Automatically "frees" anything once the program has lost all references to it
- Affects performance, but avoid memory leaks
- Java has this, C doesn't

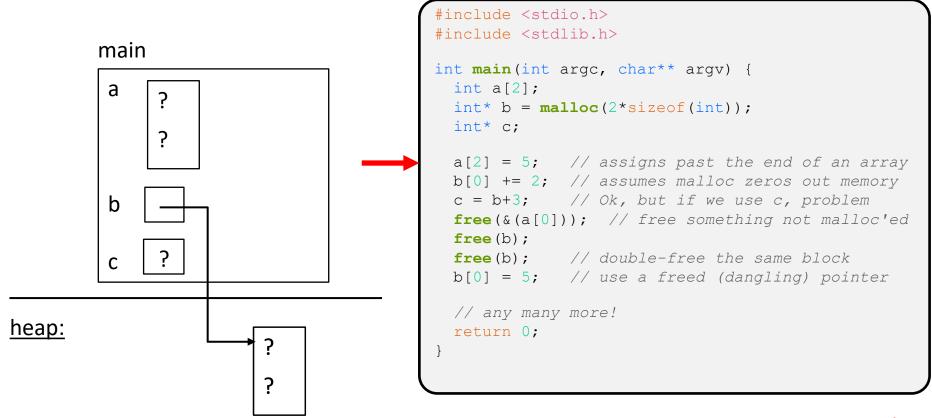


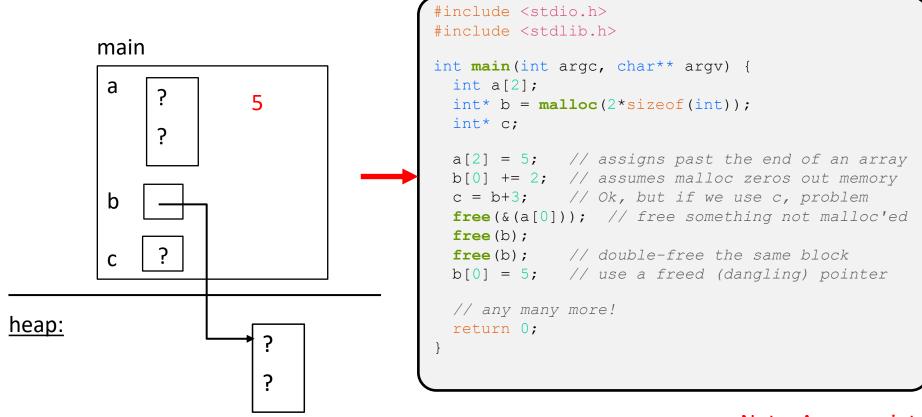
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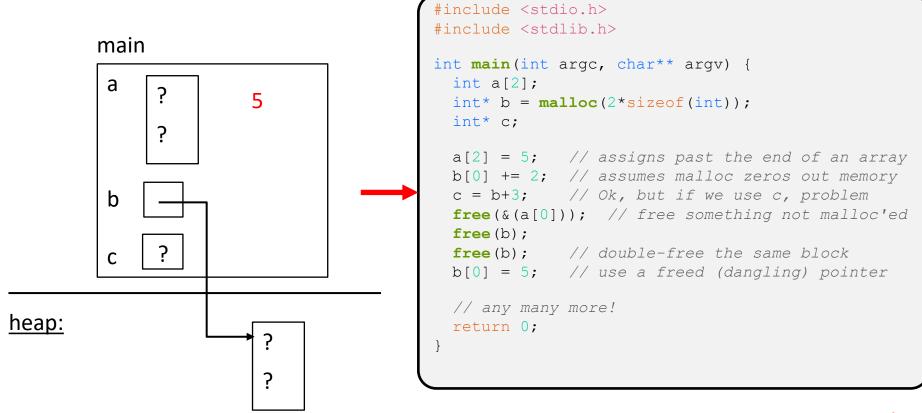
- Which line below is first to (most likely) cause a crash?
 - Yes, there are a lot of bugs, but not all cause a crash ©
 - See if you can find all the bugs!

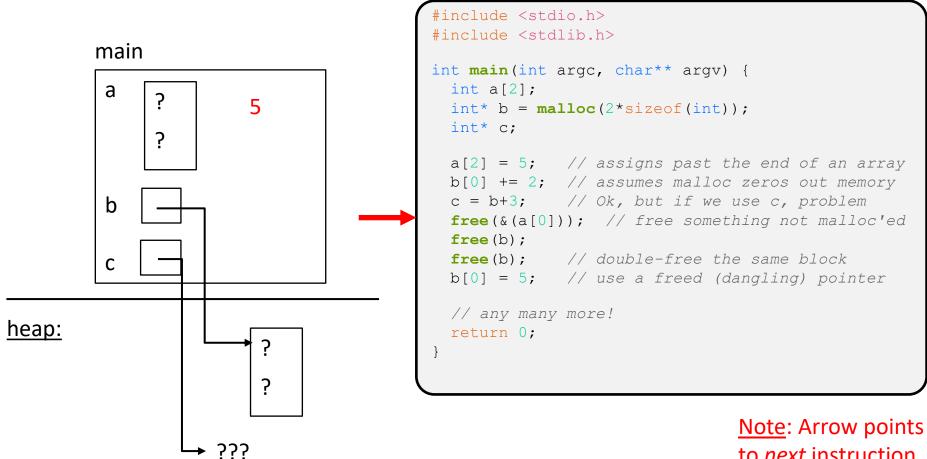
```
#include <stdio.h>
#include <stdlib.h>
int main(int argc, char** argv) {
  int a[2];
  int* b = malloc(2*sizeof(int));
  int* c;
  a[2] = 5;
 b[0] += 2;
  c = b + 3;
  free (& (a[0]));
  free(b);
  free(b);
 b[0] = 5;
  return 0;
```



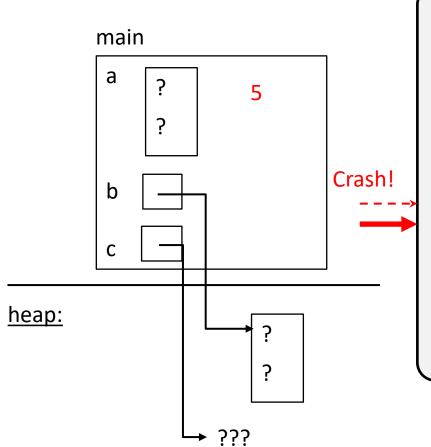




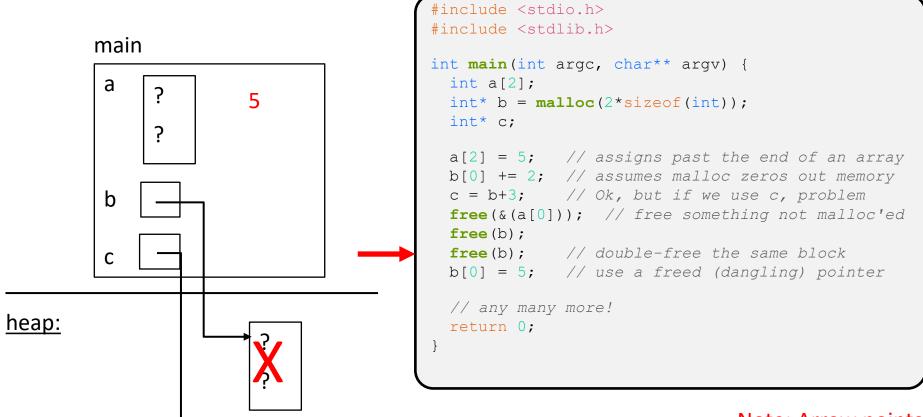




to *next* instruction.



```
#include <stdio.h>
#include <stdlib.h>
int main(int argc, char** argv) {
 int a[2];
 int* b = malloc(2*sizeof(int));
 int* c;
 a[2] = 5; // assigns past the end of an array
 b[0] += 2; // assumes malloc zeros out memory
 c = b+3; // Ok, but if we use c, problem
 free(&(a[0])); // free something not malloc'ed
 free (b);
 free(b); // double-free the same block
 b[0] = 5; // use a freed (dangling) pointer
 // any many more!
 return 0;
```



Note: Arrow points to *next* instruction.

This "double free" would also cause the program to crash

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