Recitation 04

Midterm Review!



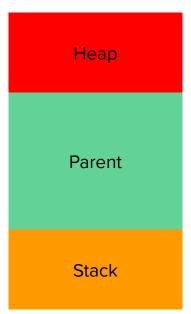
Table of Contents

- 1. Brief overview of all topics
- 2. Some thinking questions
- 3. Any questions on Midterm??
- 4. Open OH for remaining time

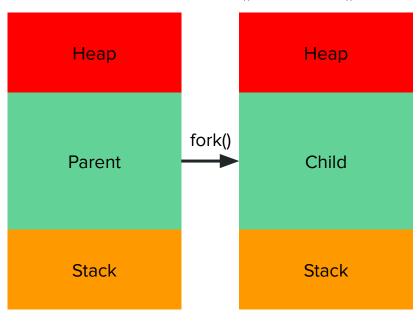
List of Midterm Topics

- 1. Processes
- 2. Signals
- 3. Terminal Control
- 4. File System
- 5. Cache, Locality and Buffering

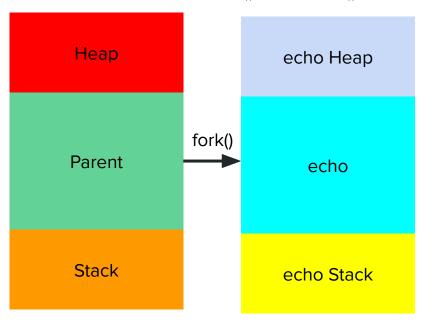
Shredder - fork(), exec(), wait()



Shredder - fork(), exec(), wait()

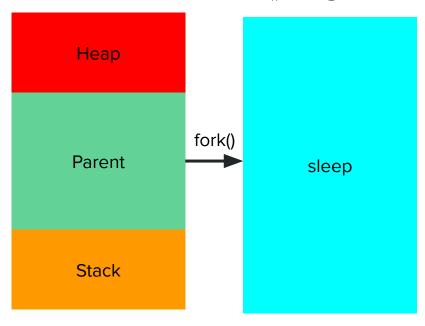


Shredder - fork(), exec(), wait()



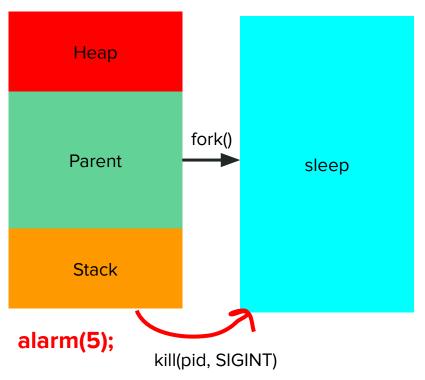
wait(pid);

shredder> echo hi hi

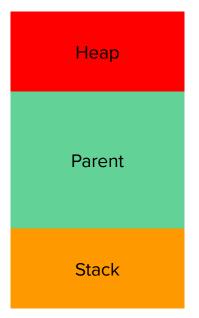


alarm(5);

shredder> echo hi
hi
shredder> sleep 10

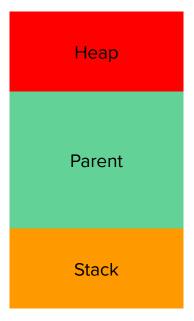


shredder> echo hi
hi
shredder> sleep 10



alarm(5); wait();

shredder> echo hi
hi
shredder> sleep 10
BWAHAHA...



alarm(5); wait(); shredder> echo hi
hi
shredder> sleep 10
BWAHAHA...
shredder>

cat < file.txt | grep comrade | wc -l -c > out.txt

penn-shell (parent)

0: stdin

1: stdout

2: stderr

3: pipe1[0]

4: pipe1[1]

5: pipe2[0]

```
cat < file.txt | grep comrade | wc -l -c > out.txt
penn-shell
 (parent)
 0: stdin
 1: stdout
                                  pipe1
 2: stderr
3: pipe1[0]
4: pipe1[1]
                                  pipe2
5: pipe2[0]
```

cat < file.txt | grep comrade | wc -l -c > out.txt

penn-shell (parent)

0: stdin

1: stdout

2: stderr

3: pipe1[0]

4: pipe1[1]

5: pipe2[0]

6: pipe2[1]



0: stdin

1: stdout

2: stderr

3: pipe1[0]

4: pipe1[1]

5: pipe2[0]

cat < file.txt | grep comrade | wc -l -c > out.txt

penn-shell (parent)

0: stdin

1: stdout

2: stderr

3: pipe1[0]

4: pipe1[1]

5: pipe2[0]

6: pipe2[1]

Child 1 for cat

0: stdin

1: stdout

2: stderr

3: pipe1[0]

4: pipe1[1]

5: pipe2[0]

6: pipe2[1]

Child 2 for grep

0: stdin

1: stdout

2: stderr

3: pipe1[0]

4: pipe1[1]

5: pipe2[0]

cat < file.txt | grep comrade | wc -1 -c > out.txt

penn-shell (parent)

0: stdin

1: stdout

2: stderr

3: pipe1[0]

4: pipe1[1]

5: pipe2[0]

6: pipe2[1]

Child 1 for cat

0: stdin

1: stdout

2: stderr

3: pipe1[0]

4: pipe1[1]

5: pipe2[0]

6: pipe2[1]

Child 2 for grep

0: stdin

1: stdout

2: stderr

3: pipe1[0]

4: pipe1[1]

5: pipe2[0]

6: pipe2[1]

Child 3 for wc

0: stdin

1: stdout

2: stderr

3: pipe1[0]

4: pipe1[1]

5: pipe2[0]

"Redirect if needed!"

Shell - pipe, redirection

cat < file.txt | grep comrade | wc -l -c > out.txt

penn-shell (parent)

0: stdin

1: stdout

2: stderr

3: pipe1[0]

4: pipe1[1]

5: pipe2[0]

6: pipe2[1]

Child 1 for cat

0: stdin

1: stdout

2: stderr

3: pipe1[0]

4: pipe1[1]

5: pipe2[0]

6: pipe2[1]

Child 2 for grep

0: stdin

1: stdout

2: stderr

3: pipe1[0]

4: pipe1[1]

5: pipe2[0]

6: pipe2[1]

Child 3 for wc

0: stdin

1: stdout

2: stderr

3: pipe1[0]

4: pipe1[1]

5: pipe2[0]

"Redirect if needed!"

Shell - pipe, redirection

cat < file.txt | grep comrade | wc -l -c > out.txt

penn-shell (parent)

0: stdin

1: stdout

2: stderr

3: pipe1[0]

4: pipe1[1]

5: pipe2[0]

6: pipe2[1]

Child 1 for cat

O: file.txt

1: stdout

2: stderr

3: pipe1[0]

4: pipe1[1]

5: pipe2[0]

6: pipe2[1]

Child 2 for grep

0: stdin

1: stdout

2: stderr

3: pipe1[0]

4: pipe1[1]

5: pipe2[0]

6: pipe2[1]

Child 3 for wc

0: stdin

1: out.txt

2: stderr

3: pipe1[0]

4: pipe1[1]

5: pipe2[0]

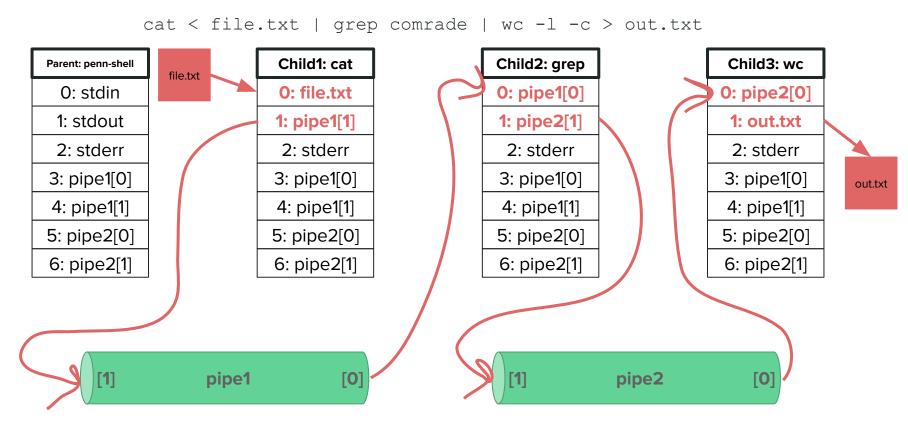
"Connect children with pipes!"

Shell - pipe, redirection

cat < file.txt | grep comrade | wc -l -c > out.txt Child1: cat Child2: grep Parent: penn-shell Child3: wc 0: stdin 0: stdin 0: file.txt 0: stdin 1: stdout 1: stdout 1: stdout 1: out.txt 2: stderr 2: stderr 2: stderr 2: stderr 3: pipe1[0] 3: pipe1[0] 3: pipe1[0] 3: pipe1[0] 4: pipe1[1] 4: pipe1[1] 4: pipe1[1] 4: pipe1[1] 5: pipe2[0] 5: pipe2[0] 5: pipe2[0] 5: pipe2[0] 6: pipe2[1] 6: pipe2[1] 6: pipe2[1] 6: pipe2[1] [0] [0] pipe1 pipe2

Content from "file.txt" will travel through pipes from Child 1 to 3 Each child will execute according to data they read from pipe/file

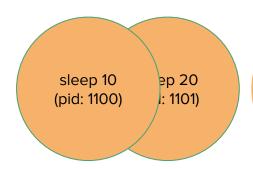
Shell - pipe, redirection





status: running **terminal control**: yes

```
shell> sleep 10 | sleep 20 &
shell> sleep 200 &
shell> sleep 300 &
shell>
```



status: running **terminal control**: no

ppid: 1000 **pgid**: 1100



status: running **terminal control**: no

ppid: 1000 **pgid**: 1150

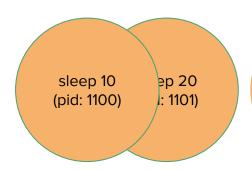


status: running **terminal control**: no

ppid: 1000 **pgid**: 1160



status: running terminal control: no



status: running **terminal control**: no

ppid: 1000 **pgid**: 1100



status: running **terminal control**: no

ppid: 1000 **pgid**: 1150



status: running
terminal control: yes

ppid: 1000 **pgid**: 1160



status: running terminal control: no

```
shell> sleep 10 | sleep 20 &
shell> sleep 200 &
shell> sleep 300 &
shell> fq
     (running sleep 300)
     (sleep 10 | sleep 20 finishes in background)
```

sleep 10 ep 20 (pid: 1100) : 1101)

> status: terminated (zombie)

terminal control: no **ppid**: 1000

pgid: 1100

sleep 200 (pid: 1150)

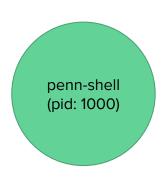
status: running terminal control: no

ppid: 1000 **pgid**: 1150

sleep 300 (pid: 1160)

status: running terminal control: yes

ppid: 1000 **pgid**: 1160



status: running terminal control: yes

```
shell> sleep 10 | sleep 20 &
shell> sleep 200 &
shell> sleep 300 &
shell> fq
     (running sleep 300)
     (sleep 10 | sleep 20 finishes in background)
^7.
```

Finished: sleep 10 | sleep 20 shell>

```
sleep 10
                ep 20
(pid: 1100)
                 : 1101)
```

status: terminated (zombie) terminal control: no

ppid: 1000 **pgid**: 1100

sleep 200 (pid: 1150)

status: running terminal control: no

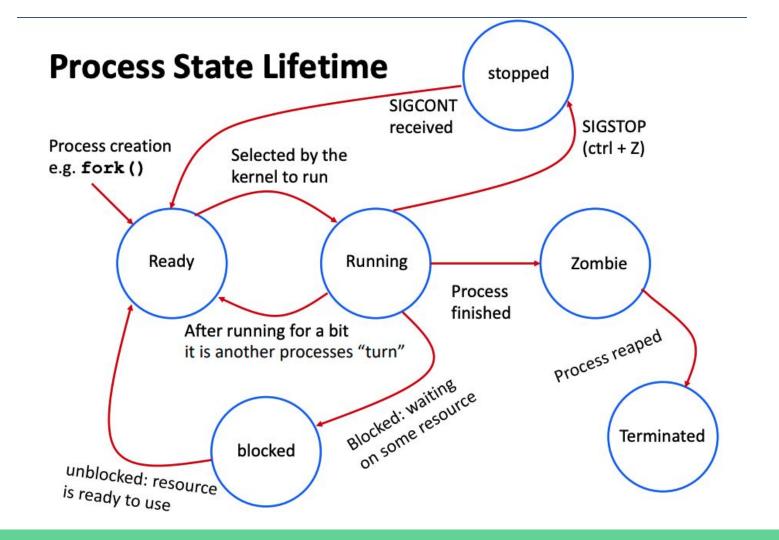
ppid: 1000 **pgid**: 1150

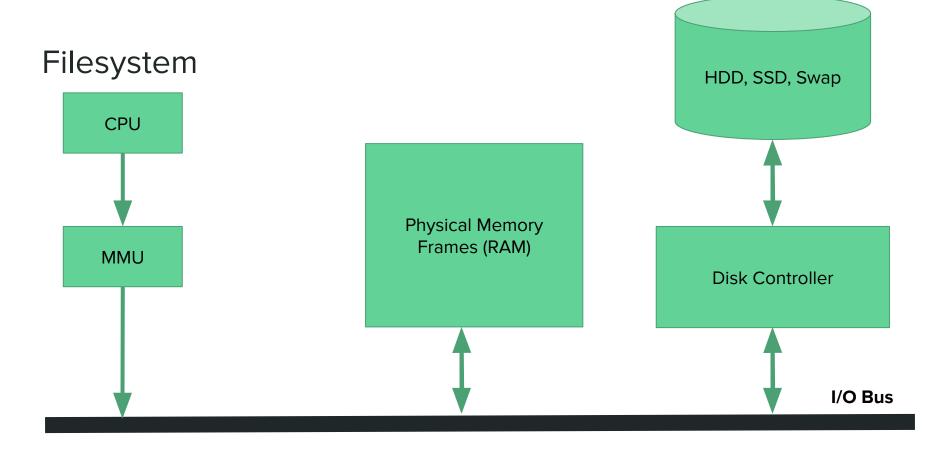
sleep 300 (pid: 1160)

status: stopped terminal control: no

ppid: 1000 **pgid**: 1160

```
(SIGTSTP sent to foreground process, shells re-gains TC)
(Any terminated child in background is waited on)
(Reprompt)
```





Filesystem

- Interfaces, abstractions to organize/access data in disk
- Sequence of bytes ⇒ Blocks of bytes ⇒ Files
- How and where to access?
 - Directory file
 - o Bitmap
 - FAT
 - Inodes

Filesystem - Some questions

- What are the pros and cons of each method?
- Is bitmap alone enough for a filesystem API? If not, how can we modify?
- Contiguous allocation vs Linked List allocation
- Internal Fragmentation and External Fragmentation
- Difference between Inodes and FAT?

Buffering

```
int main() {
  for (i is [1,3]) {
    pid = fork();
    if (pid == 0) print("hi");
  }
  i=2
}
```

i=3

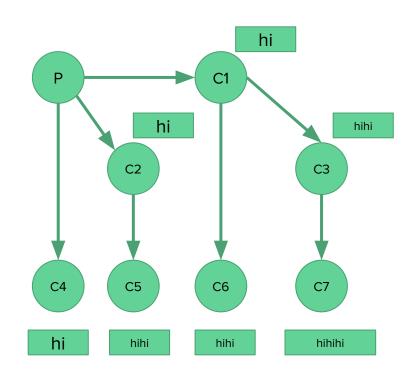
How many children?

1

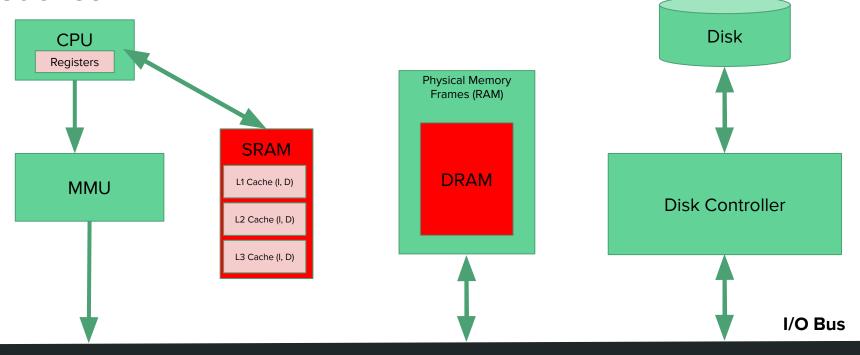
How many "hi\n"?

12

How to resolve? Why do we buffer?



Caches



Least Recently Used (LRU)

- Memory is limited
- Which line of memory to evict/replace when we run out of memory?
 - Least Recently Used
- Advantages of LRU
 - Generally good performance, we are evicting a page that is "least" frequently used
 - Reduces number of page faults
- Disadvantages of LRU
 - Quite costly to find the LRU page
- Think of a scenario where LRU may actually hurt performance
 - Sequential Access: If some sequential access pattern forces LRU to evict and re-allocate parts of memory, we will have poor performance
- How else can we design the eviction policy?

Some Thinking Questions 1: What might go wrong here?

```
void poll() {
         for (all background jobs) {
              waitpid(-pgid, status, WNOHANG |
   WUNTRACED);
              modify queue();
   while (true) {
10
        poll();
        prompt and get command();
12
        pipe();
        for (all children) fork(), setpgid();
14
              pipe(), redirect(), exec()
15
        if (background) { waitpid(-1, status, WNOHANG); }
         else { waitpid(-1, status, WUNTRACED); }
16
```

- 1. We are waiting for only 1 child process in a job in lines 13, 14
 - a. What if "sleep 10 | sleep 100"?
- 2. We are waiting for "any child process" waitpid(-1) in line 13
 - a. What if there were zombied processes in the background?

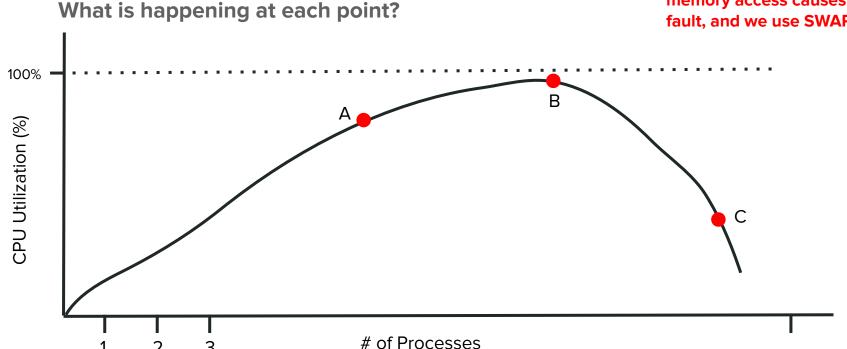
Some Thinking Questions 2

Consider this graph of CPU Utilization vs # of Processes Running

A: Context Switching becomes a problem

B: SUM(memory utilization) > RAM, so we start using more SWAP file

C: Thrashing: Most of the memory access causes a page fault, and we use SWAP a lot



Some Thinking Questions 3

Suppose our cache block size is 4096 bytes and we had a data structure that is 128 bytes big. If we had an array of this structure, what principle could we take advantage of when we fetch this array sequentially?

Spatial locality

What if the data struct is 4096 bytes big?

Cache is almost meaningless since only one data struct can fit in it

More Practice Problems

https://www.seas.upenn.edu/~cis3800/24sp/exams/midterm

Any Questions?