ESE3700: Circuit-Level Modeling, Design, and Optimization for Digital Systems

Lec 17: April 7, 2025

Memory Overview: RAM Core

Today

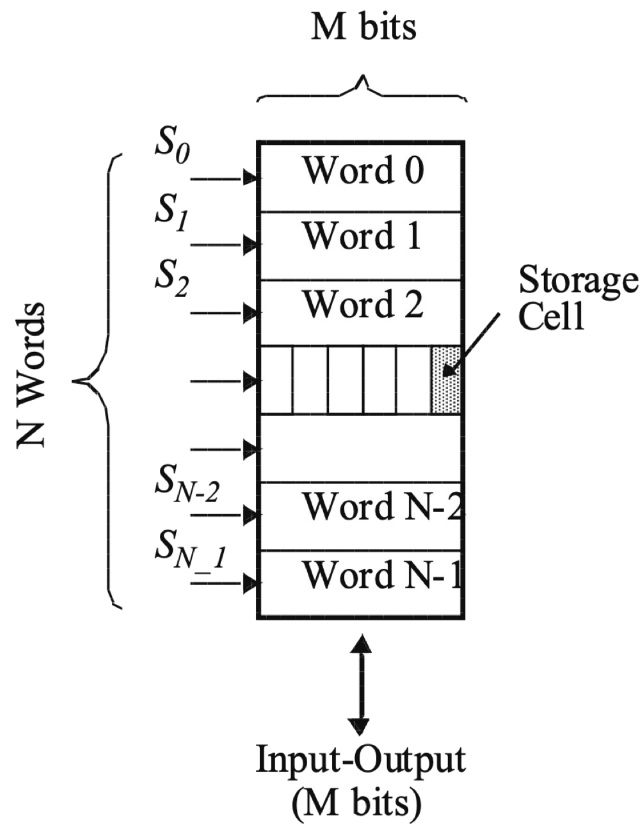
- Memory Overview
- Memory
  - Memory core
    - 6T SRAM
- Project 2 is on this

Memory Overview

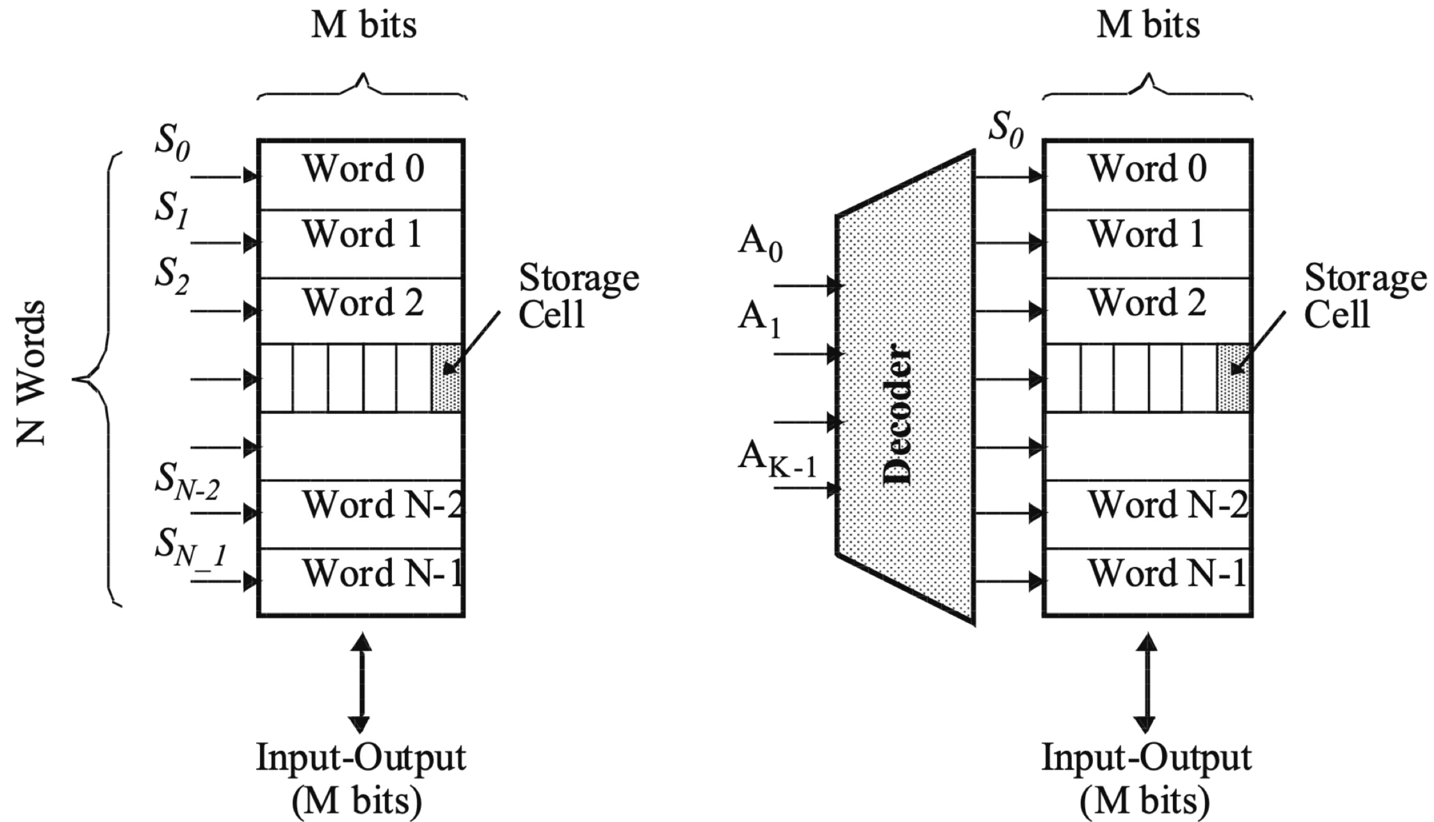
Semiconductor Memory Classification

RWM		**NVRWM**	**ROM**
Random Access	**Non-Random Access**	EPROM E^2PROM FLASH	Mask-Programmed Programmable (PROM)
SRAM DRAM	FIFO LIFO Shift Register CAM		

Memory Architecture: Core



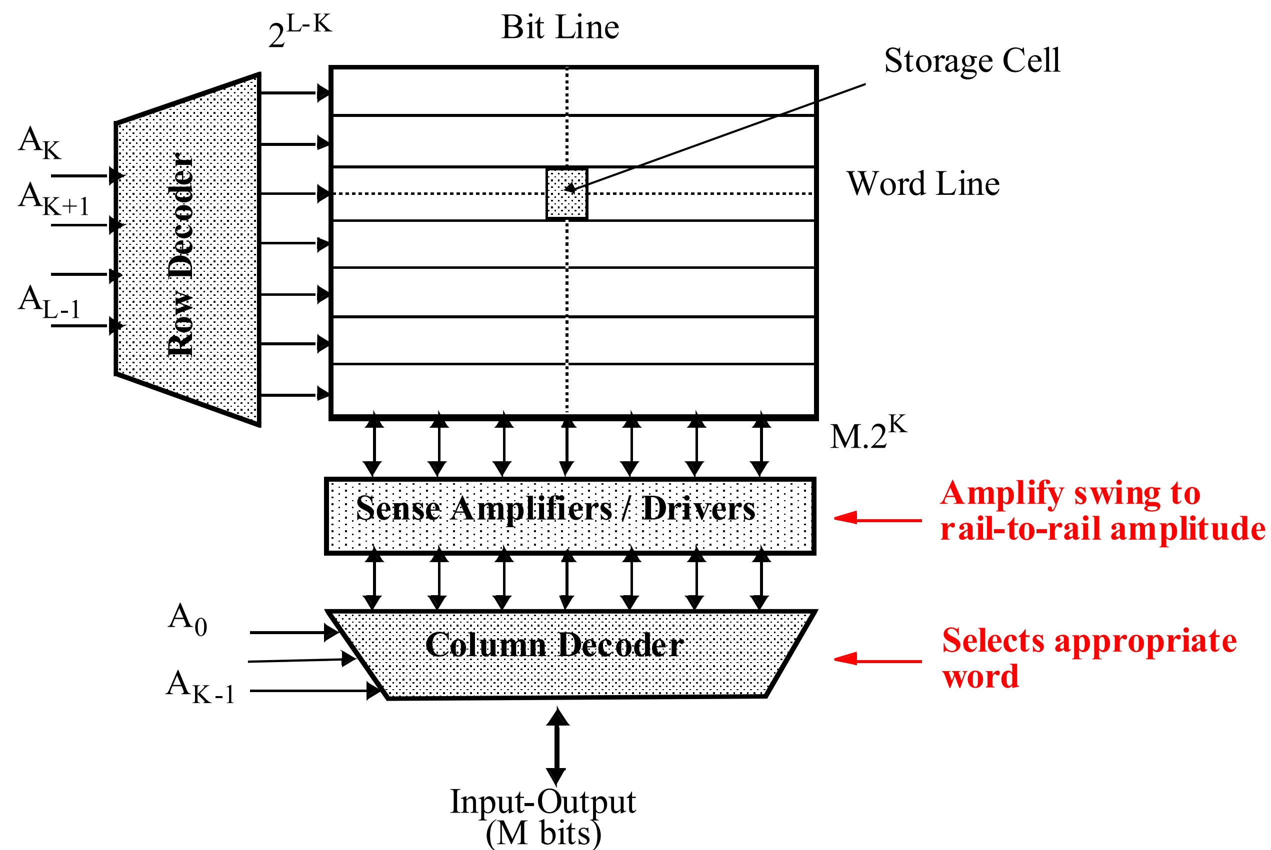
Memory Architecture: Decoders



N words => N select signals
Too many select signals

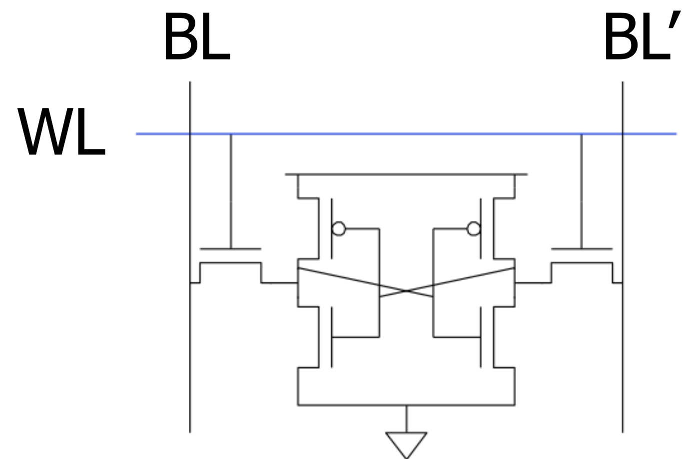
Decoder reduces # of select signals
$K = log_2 N$

Array-Structured Memory Architecture

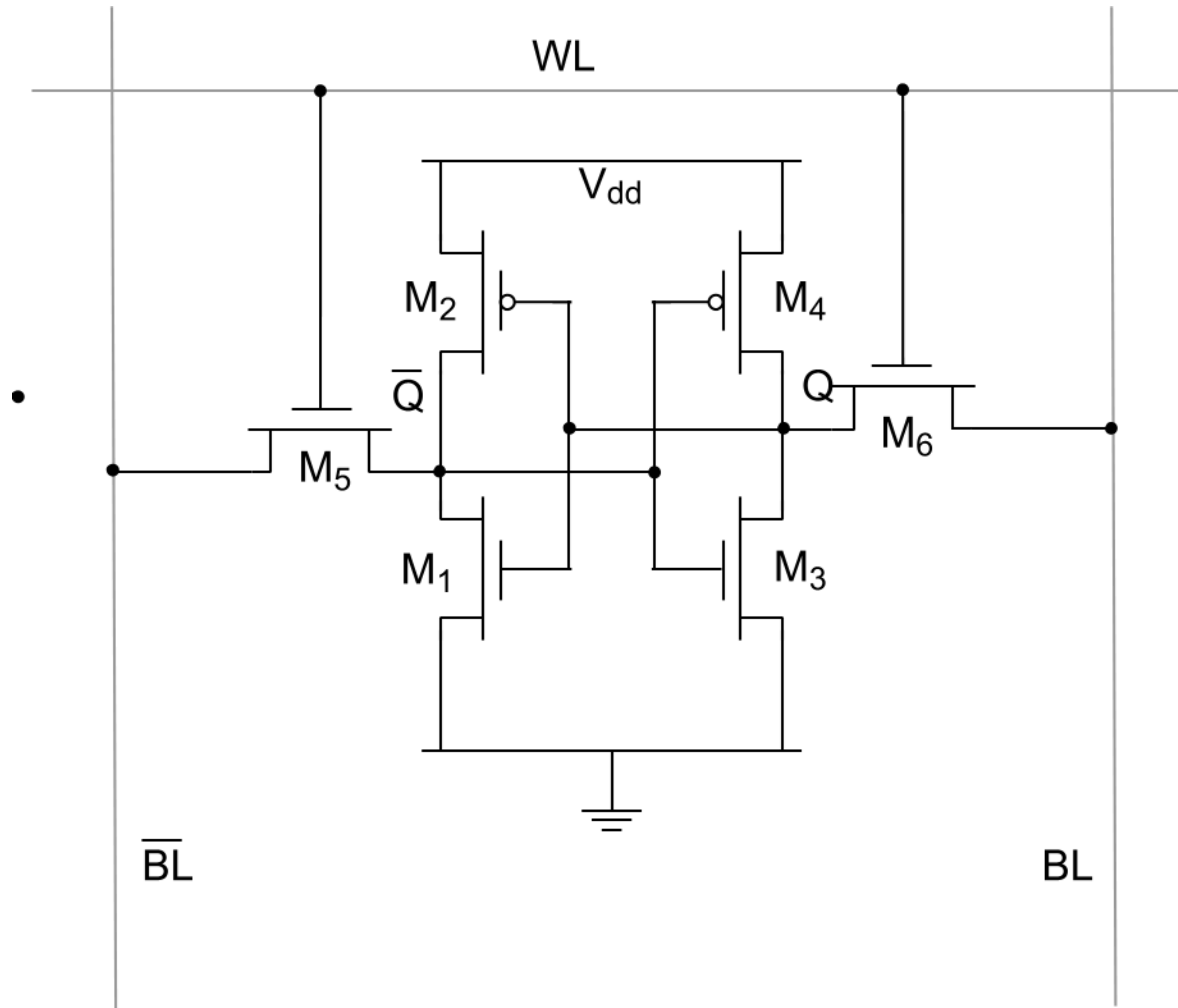


6T SRAM Cell

- Cell size accounts for most of array size
  - Reduce cell size at expense of complexity
- 6T SRAM Cell
  - Used in most commercial chips
  - Data stored in cross-coupled inverters
- Read:
  - Precharge BL, BL'
  - Raise WL
- Write:
  - Drive data onto BL, BL'
  - Raise WL



6-transistor CMOS SRAM Cell



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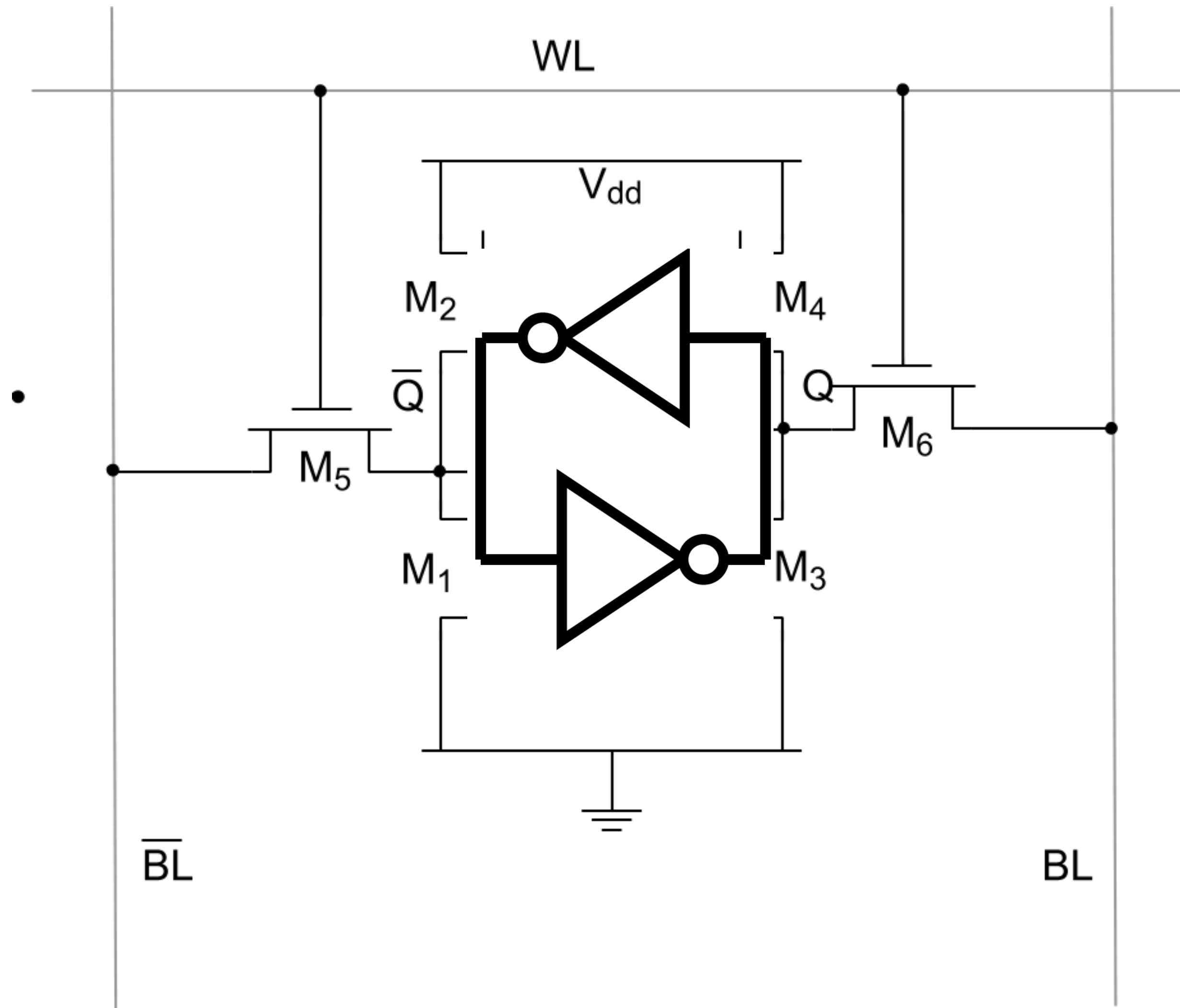
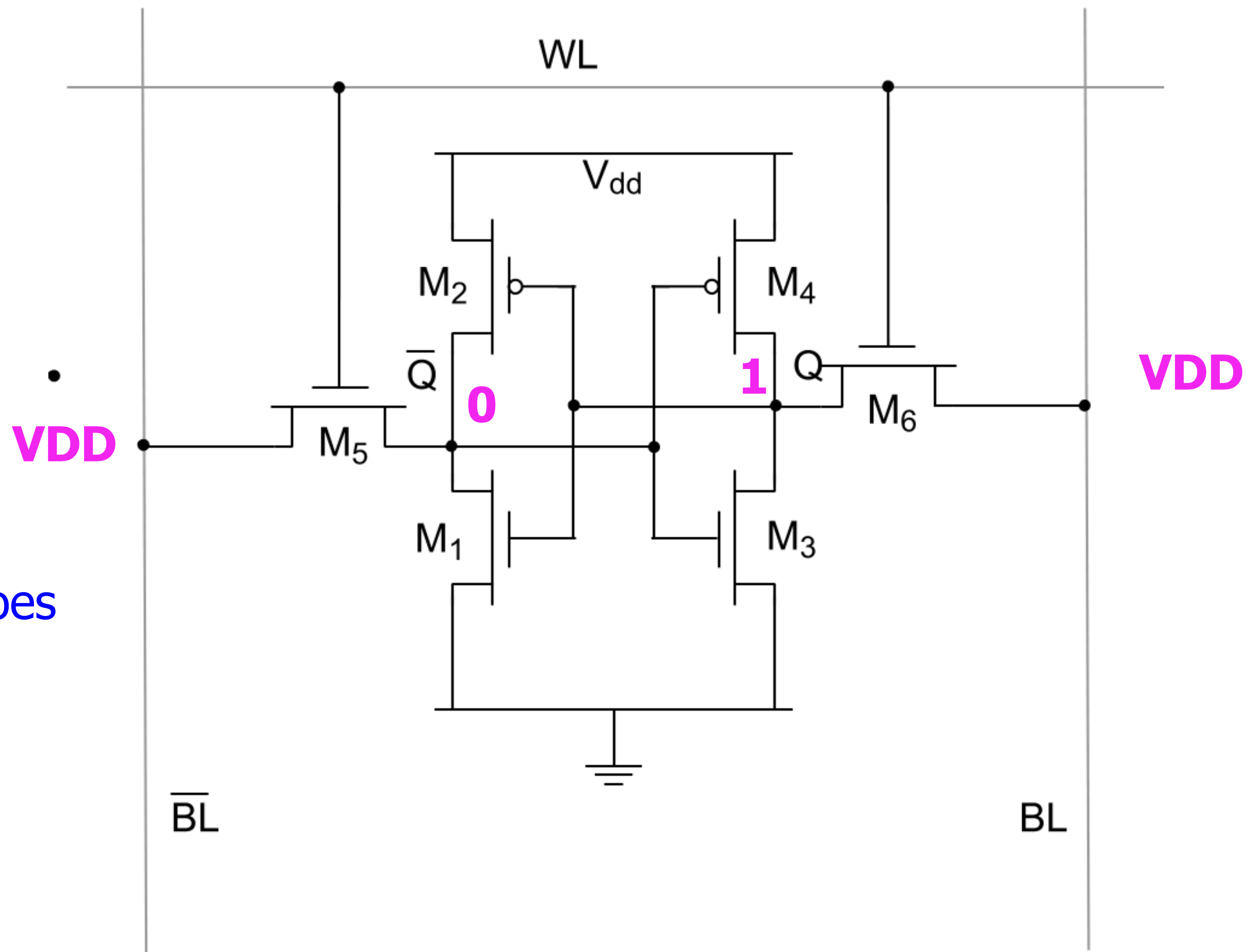


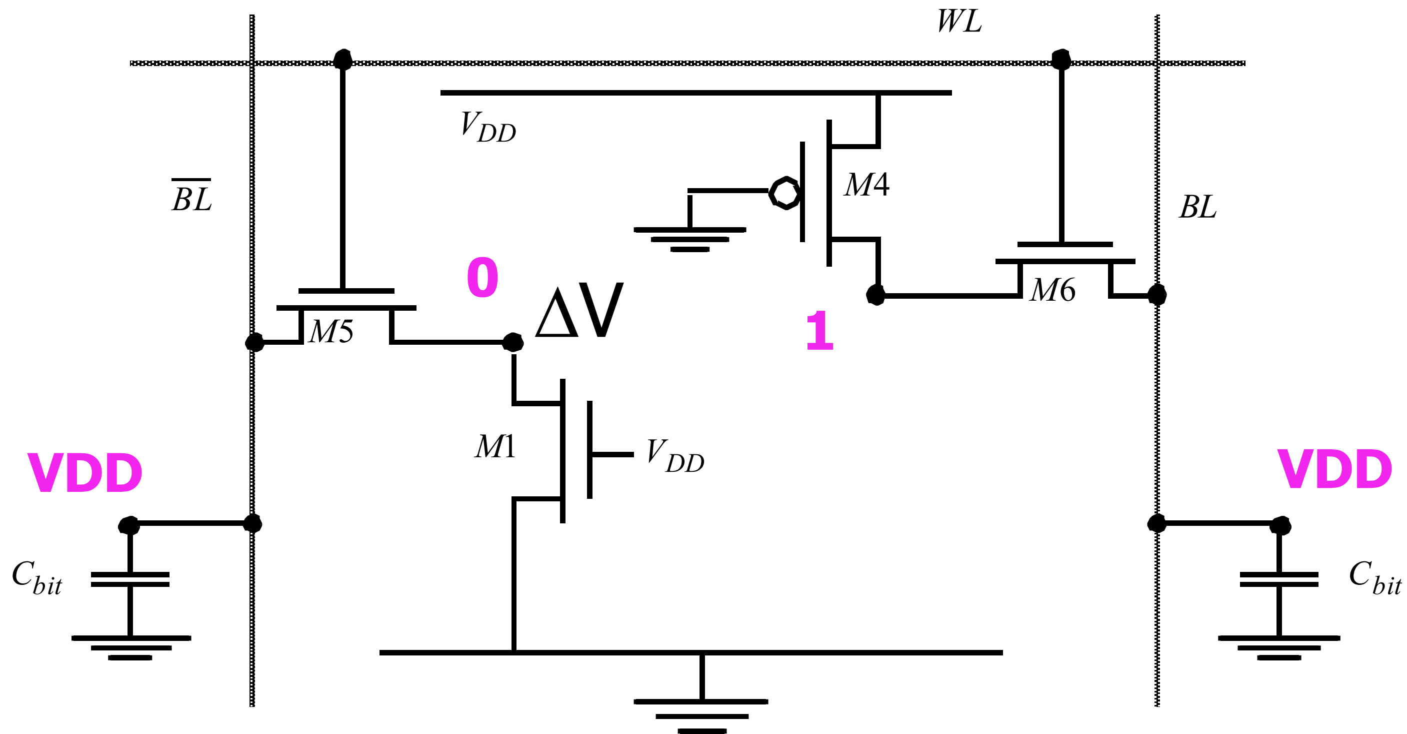
6-transistor CMOS SRAM Cell (preclass 1)

Assume 1 is stored (Q=1)

Read Operation:
- First bitlines get precharged high (Vdd)
- Then wordline goes high (Vdd)
  - Precharge disconnected

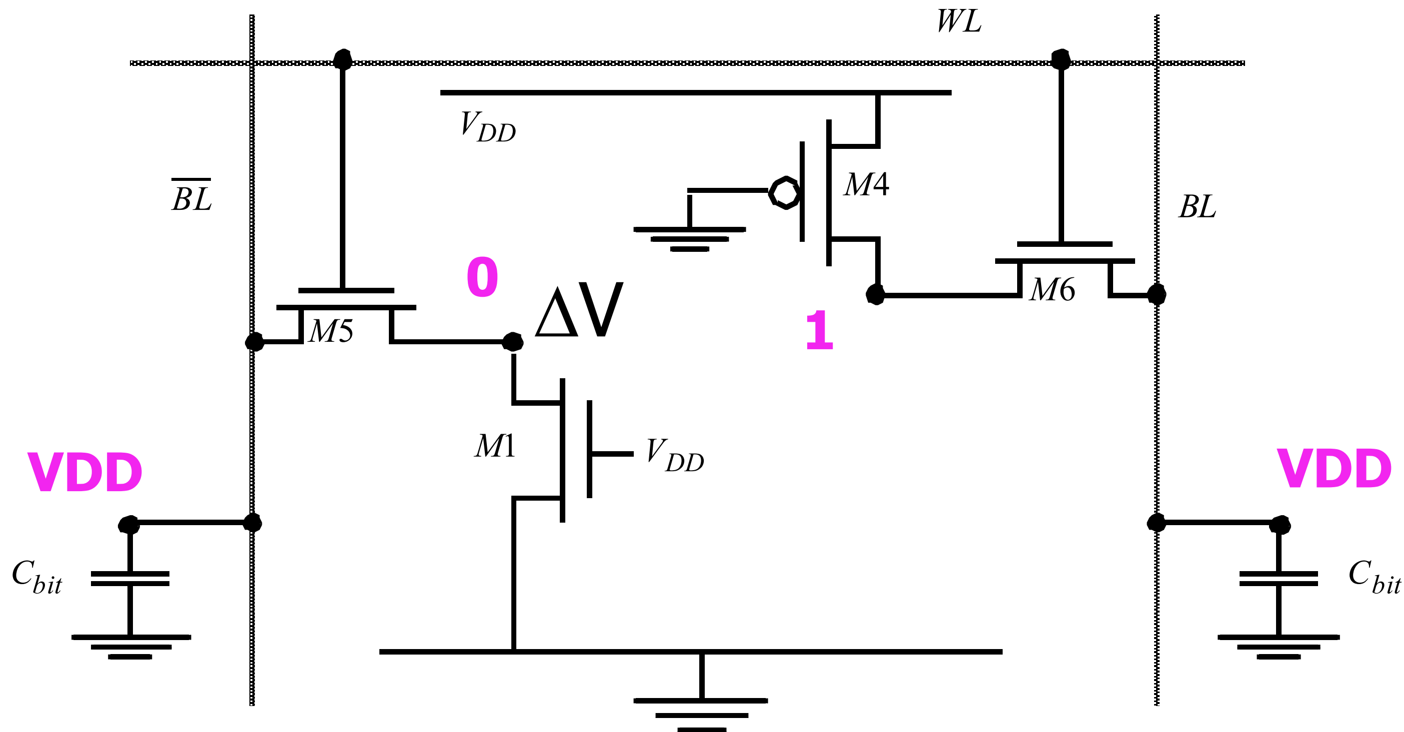


CMOS SRAM Analysis (Read) (preclass 1)



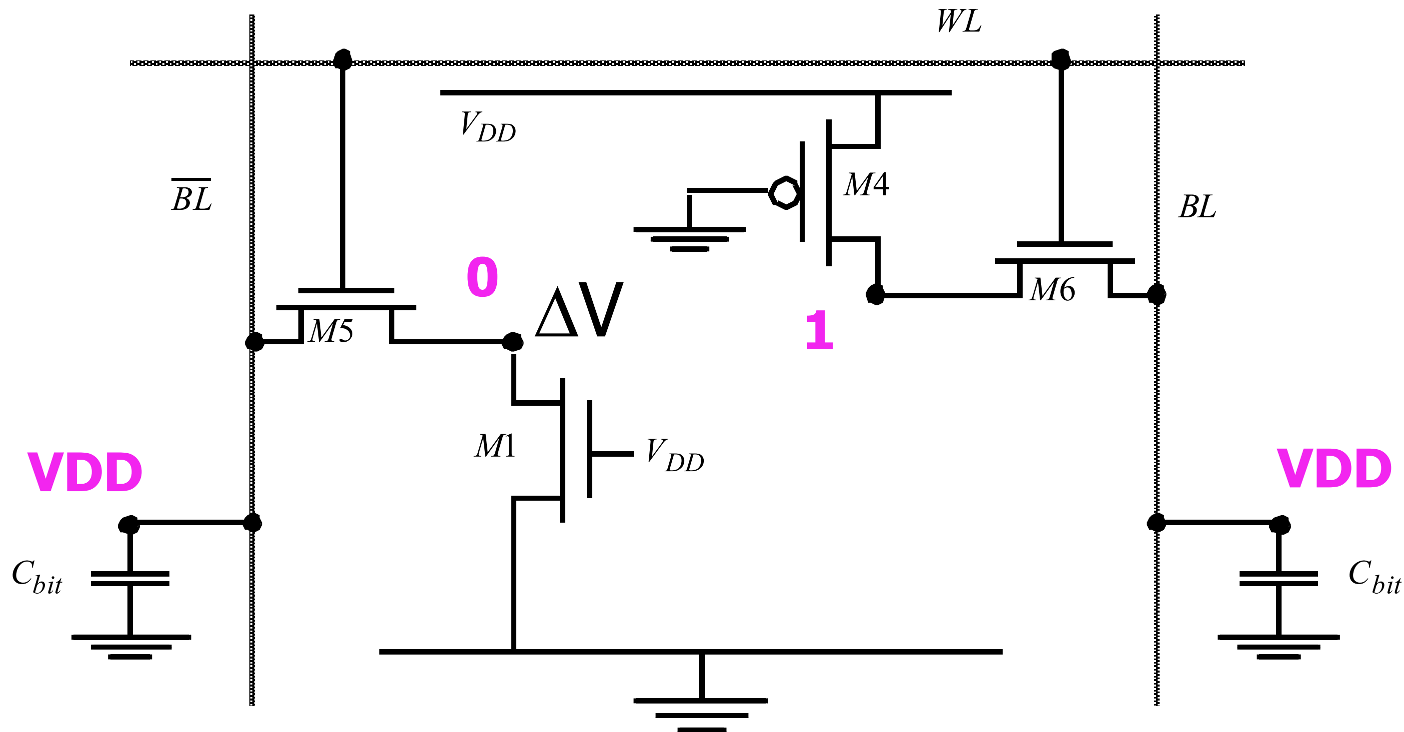
Which two transistors are discharging \BL to Gnd?
What regions of operation are the transistors in?
Write the KCL equation for the two transistors

CMOS SRAM Analysis (Read)



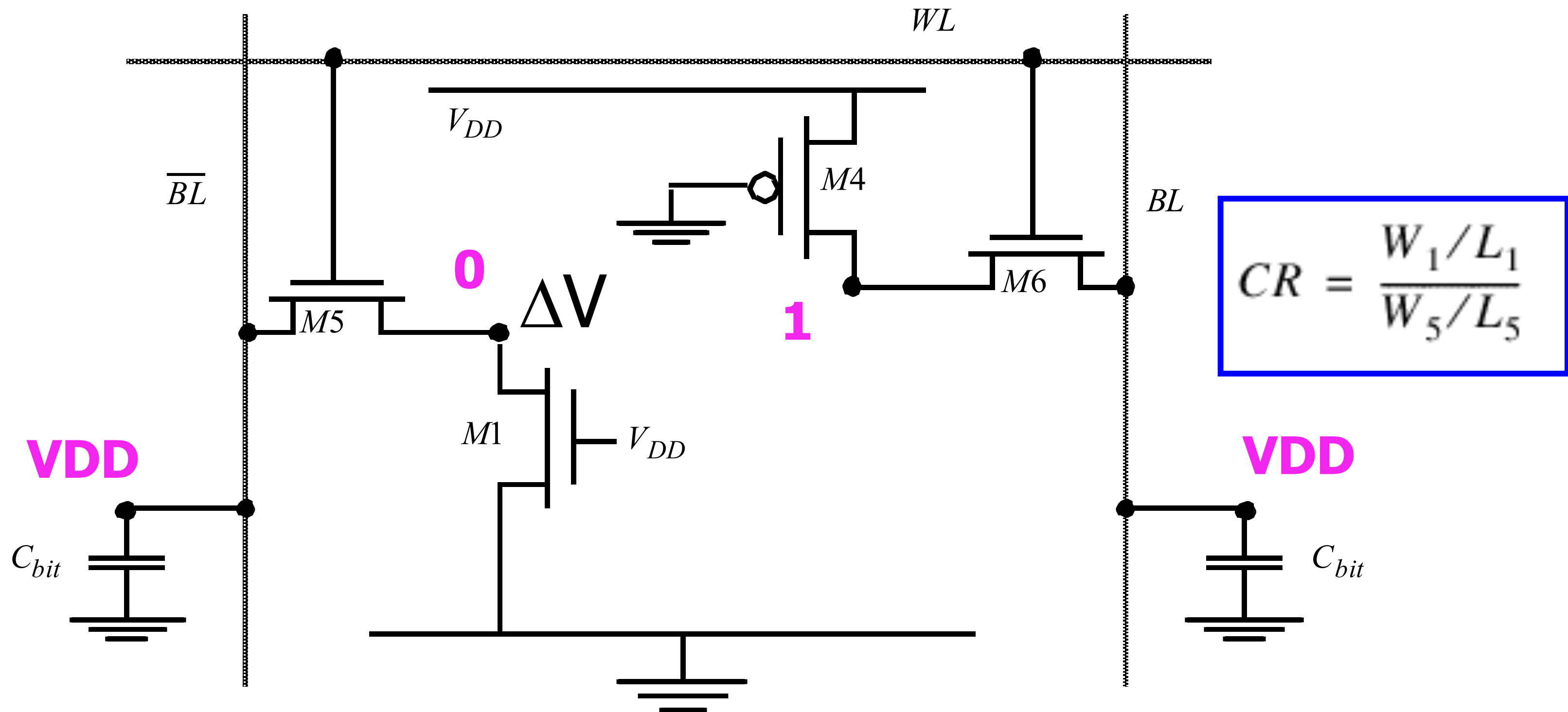
$$k_{n,M5}\left(V_{DD} - \Delta V - V_{Tn}\right)^2 = k_{n,M1}\left((V_{DD} - V_{Tn})\Delta V - \frac{\Delta V^2}{2}\right)$$

CMOS SRAM Analysis (Read)



$$\frac{k_{n,M1}}{k_{n,M5}} = \frac{\left(\frac{W}{L}\right)_1}{\left(\frac{W}{L}\right)_5} = \frac{\left(V_{DD} - \Delta V - V_{Tn}\right)^2}{(V_{DD} - V_{Tn})\Delta V - \frac{\Delta V^2}{2}}$$

CMOS SRAM Analysis (Read)

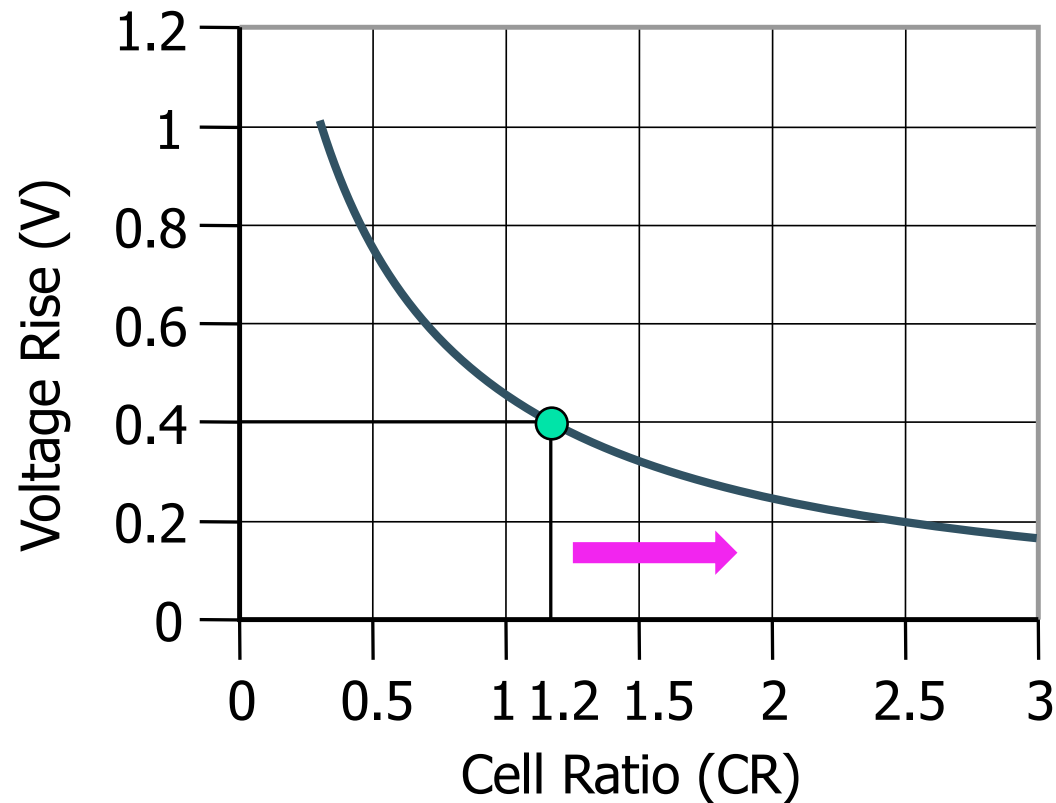


$$\frac{k_{n,M1}}{k_{n,M5}} = \frac{\left(\frac{W}{L}\right)_1}{\left(\frac{W}{L}\right)_5} = \frac{\left(V_{DD} - \Delta V - V_{Tn}\right)^2}{(V_{DD} - V_{Tn})\Delta V - \frac{\Delta V^2}{2}}$$

$\Delta V = V_{Tn}$

$$\frac{\left(\frac{W}{L}\right)_1}{\left(\frac{W}{L}\right)_5} = \frac{\left(V_{DD} - 2V_{Tn}\right)^2}{(V_{DD} - 1.5V_{Tn})V_{Tn}}$$

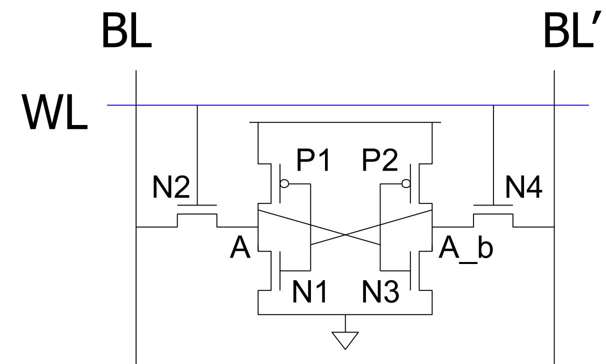
CMOS SRAM Analysis (Read)

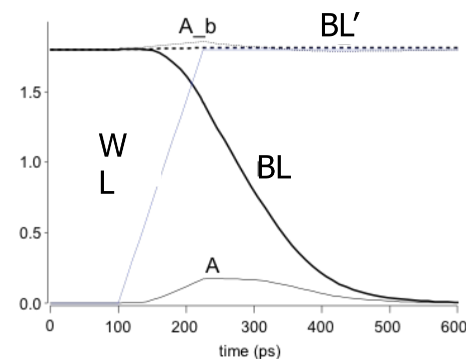


$$CR = \frac{W_1/L_1}{W_5/L_5}$$

SRAM Read

- Precharge both bitlines high
- Then turn on wordline, WL
- One of the two bitlines will be pulled down by the cell
- Ex: A = 0, A_b = 1
  - BL discharges, BL' stays high
  - But A bumps up slightly
- *Read stability*
  - A must not flip
  - N1 > N2

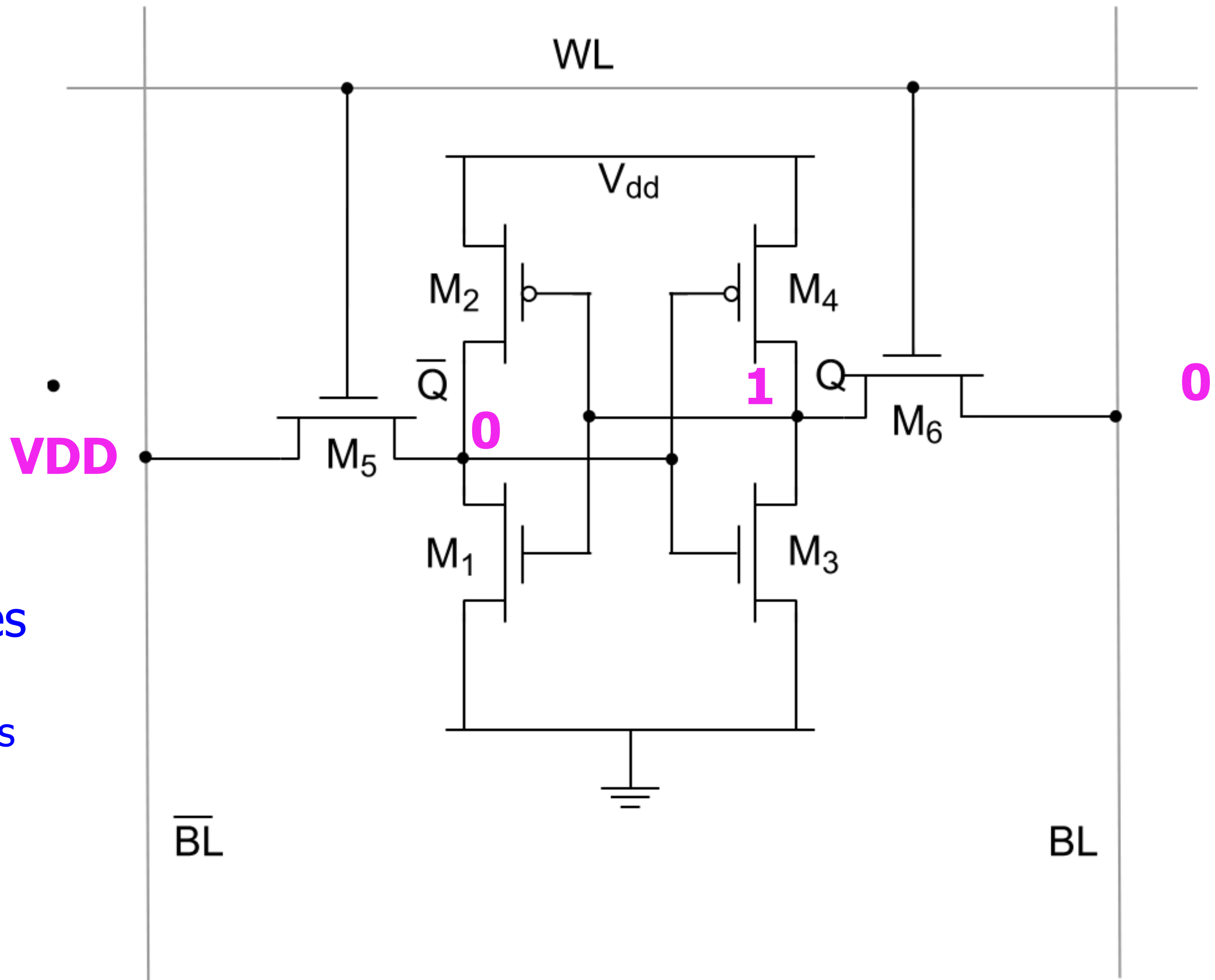




6-transistor CMOS SRAM Cell (preclass 2)

Assume 1 is stored (Q=1)

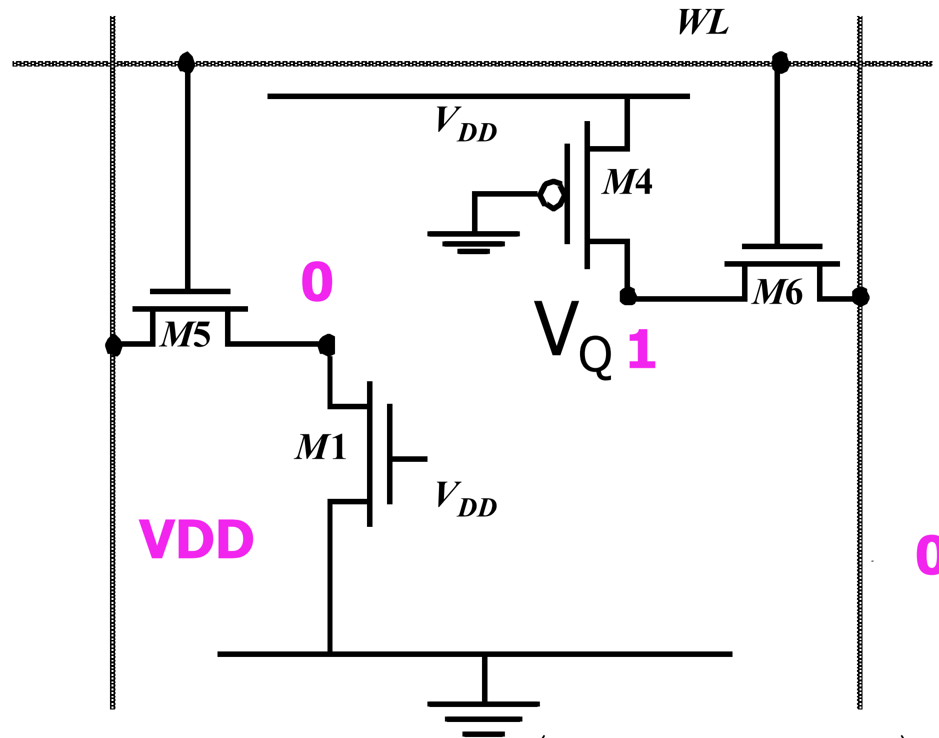
Write Operation:
- Want to write a 0
- First drive bitlines with input data
- Then wordline goes high (Vdd)
  - Still driving bitlines



CMOS SRAM Analysis (Write) (preclass 2)

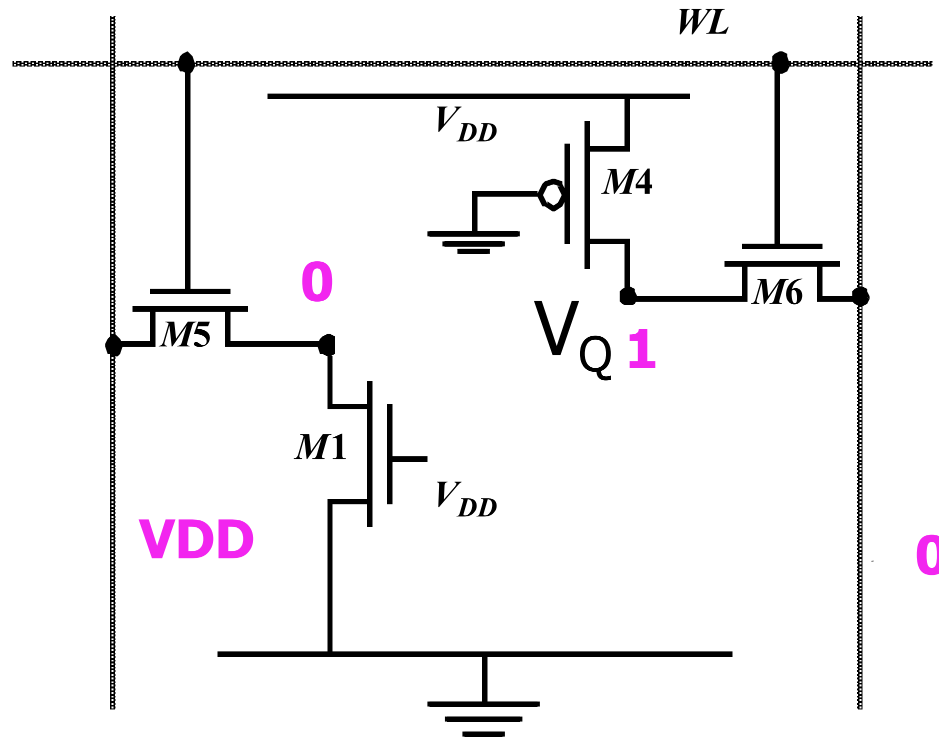
Which transistor discharges Q to Gnd?
Which transistor must it overpower?
What regions of operation are the transistors in?
Write the KCL equation for the two transistors

CMOS SRAM Analysis (Write)



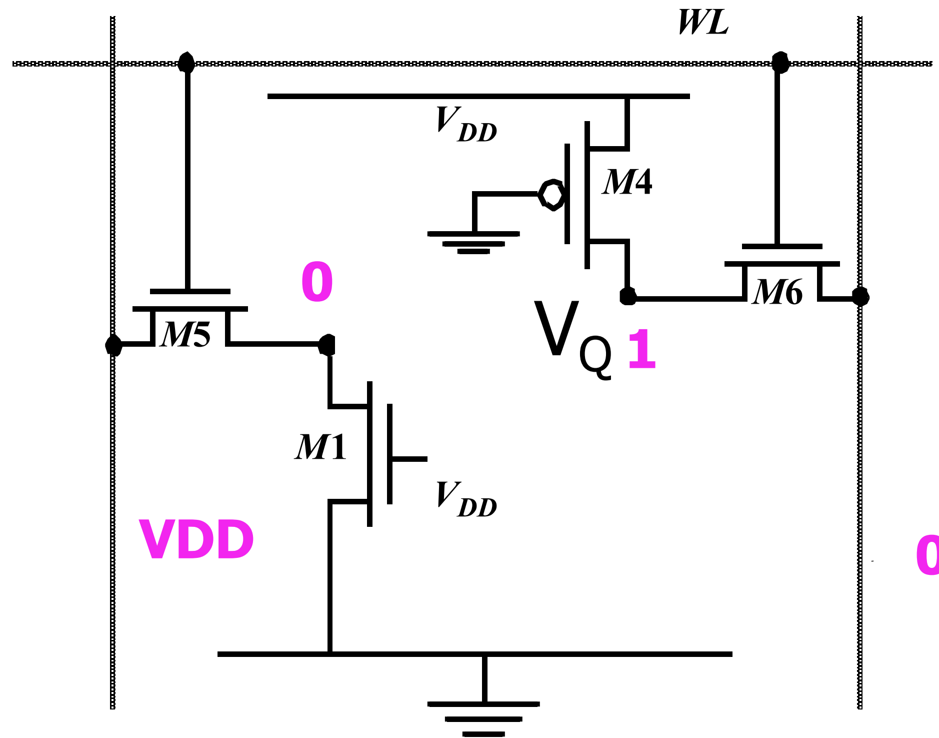
$$k_{p,M4}\left(V_{DD}-\left|V_{Tp}\right|\right)^2 = k_{n,M6}\left((V_{DD}-V_{Tn})V_Q-\frac{V_Q^2}{2}\right)$$

CMOS SRAM Analysis (Write)



$$\frac{k_{p,M4}}{k_{n,M6}} = \frac{(V_{DD} - V_{Tn})V_Q - \frac{V_Q^2}{2}}{\left(V_{DD} - \left|V_{Tp}\right|\right)^2}$$

CMOS SRAM Analysis (Write)



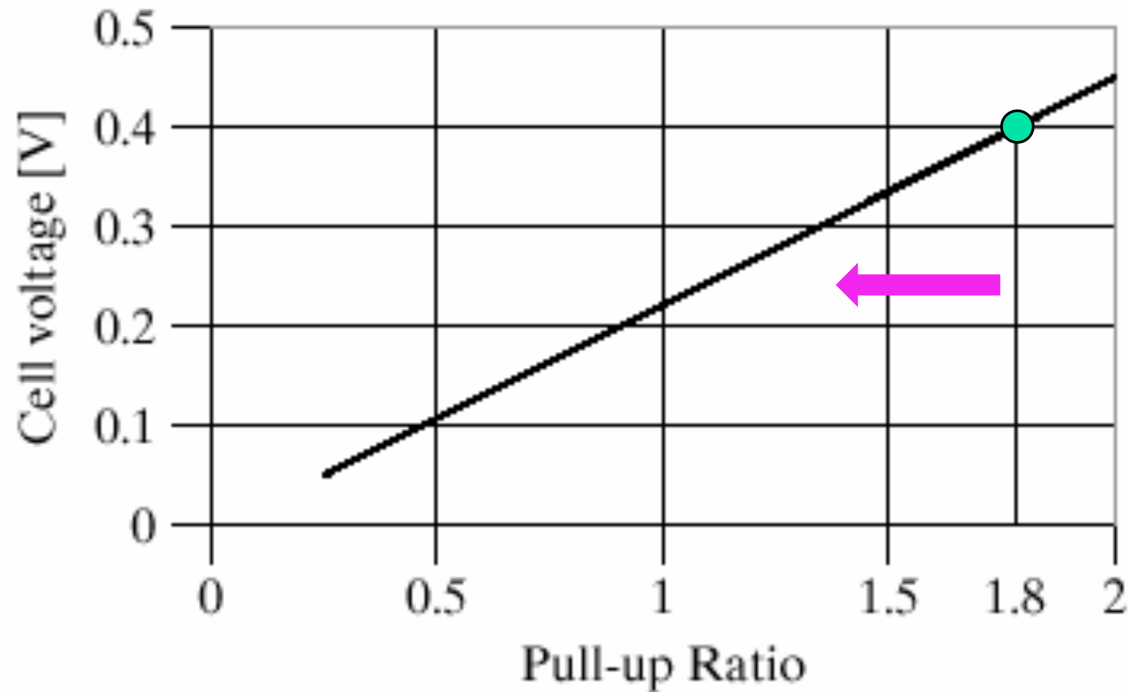
$$PR = \frac{W_4/L_4}{W_6/L_6}$$

$$\frac{k_{p,M4}}{k_{n,M6}} = \frac{(V_{DD} - V_{Tn})V_Q - \frac{V_Q^2}{2}}{\left(V_{DD} - |V_{Tp}|\right)^2}$$

$V_Q = V_{Tn}$ →

$$\frac{k_{p,M4}}{k_{n,M6}} = \frac{(V_{DD} - V_{Tn})V_{Tn} - \frac{V_{Tn}^2}{2}}{\left(V_{DD} - |V_{Tp}|\right)^2}$$

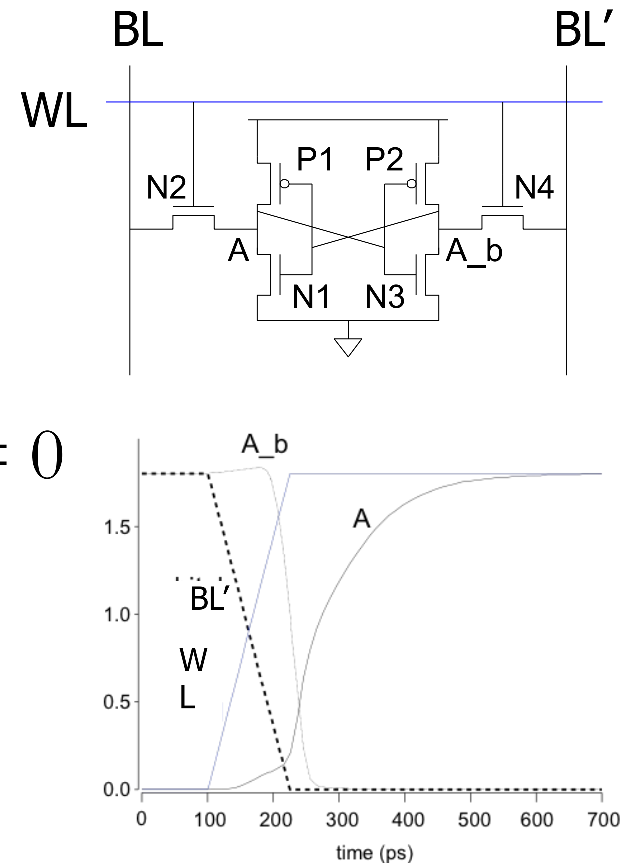
CMOS SRAM Analysis (Write)



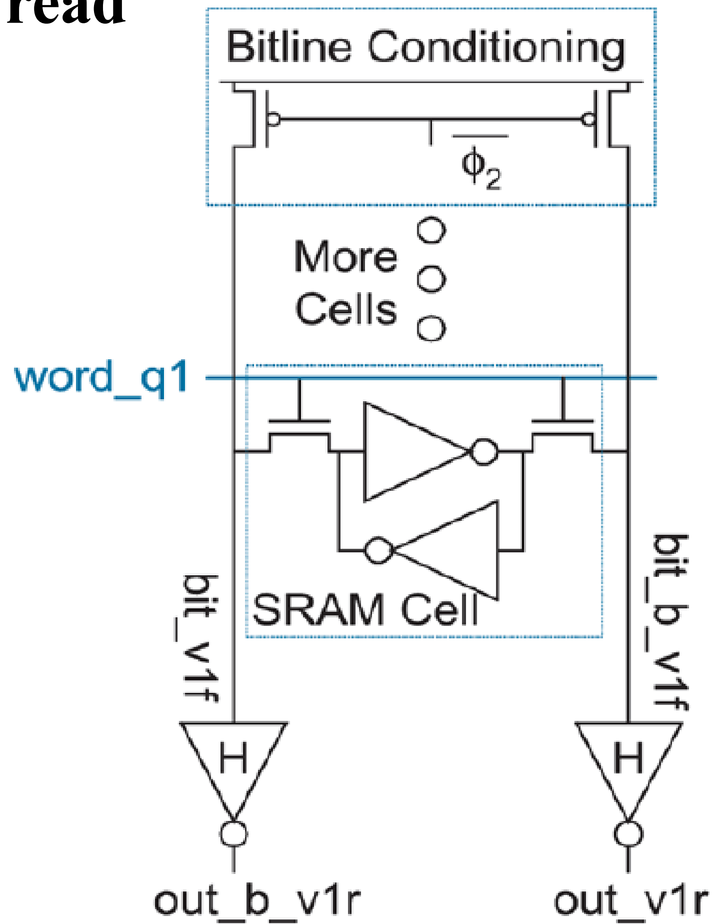
$$PR = \frac{W_4/L_4}{W_6/L_6}$$

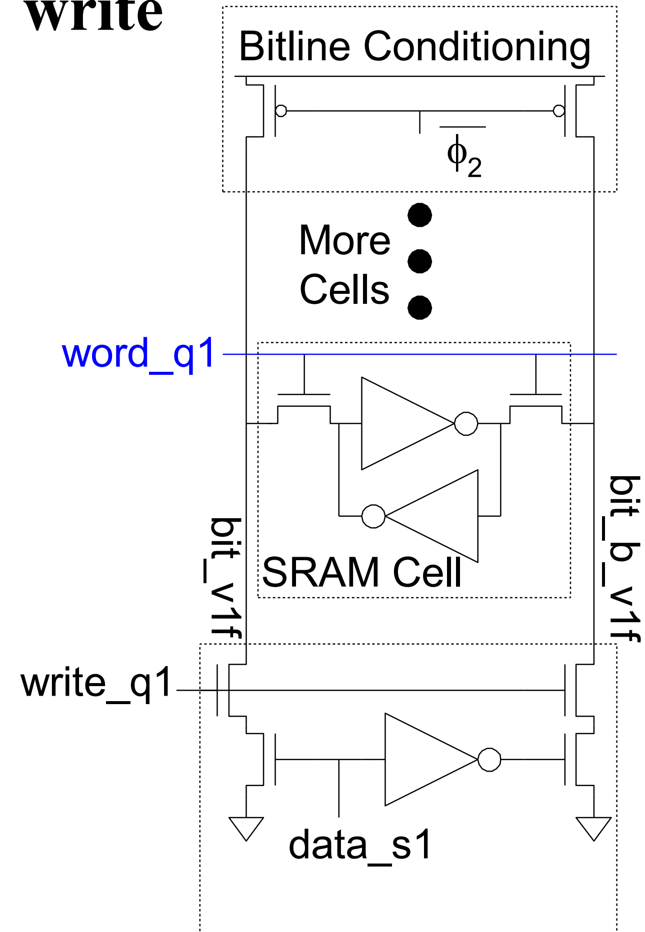
SRAM Write

- Drive one bitline high, the other low
  - Depending on write data
- Then turn on wordline, WL
- Bitlines overpower cell with new value
- Ex: A = 0, A_b = 1, BL = 1, BL' = 0
  - Force A_b low, then A charges high
- *Writability*
  - Must overpower feedback inverter
  - N4 >> P2



SRAM Column Example



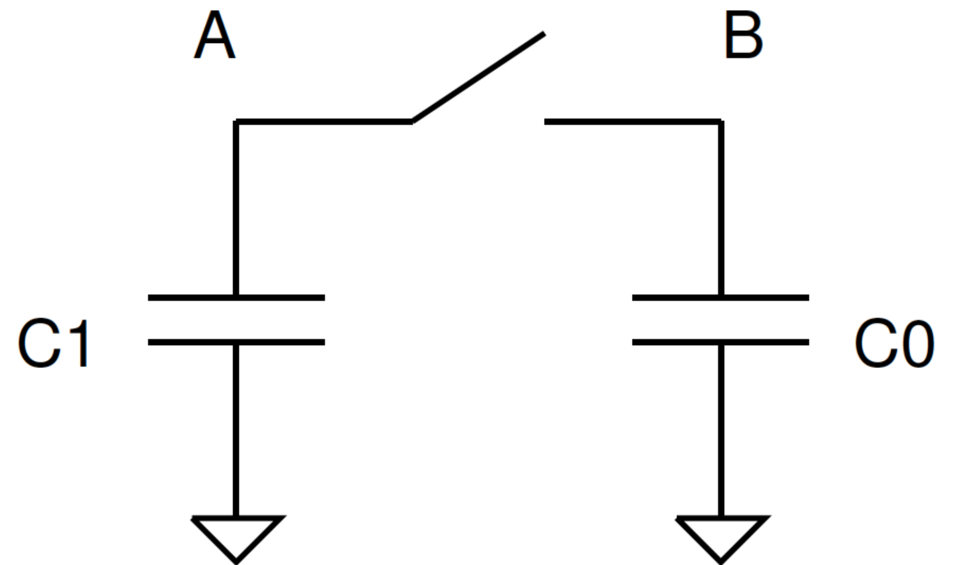


Charge Sharing

Initially

- A @ 1V
- B @ 0V

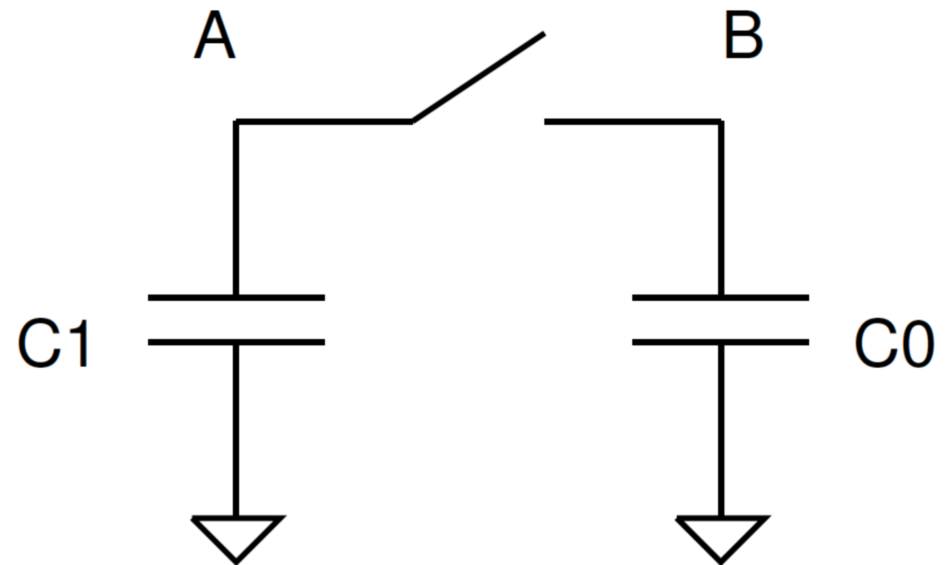
- Q_A=1V*C1=C1



Charge Sharing

Initially
- A @ 1V
- B @ 0V
- Q_A=1V*C1=C1

Close switch
- Q_{tot}=V_{final}*(C1+C0)
- Charge conservation
  - Q_A=Q_{tot}
- C1=V_{final}*(C1+C0)

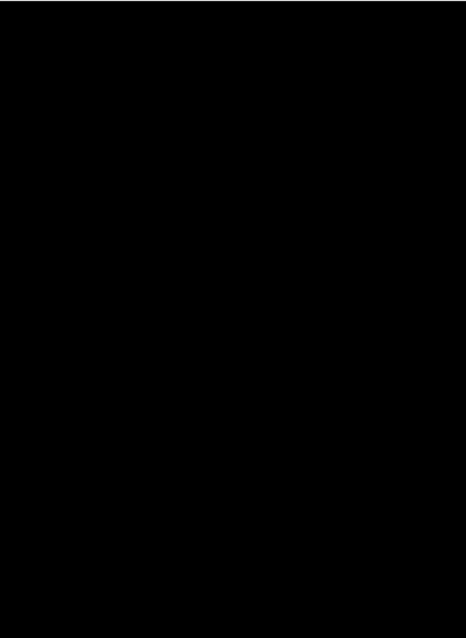


$$V_{final} = \frac{C1}{C1 + C0}$$

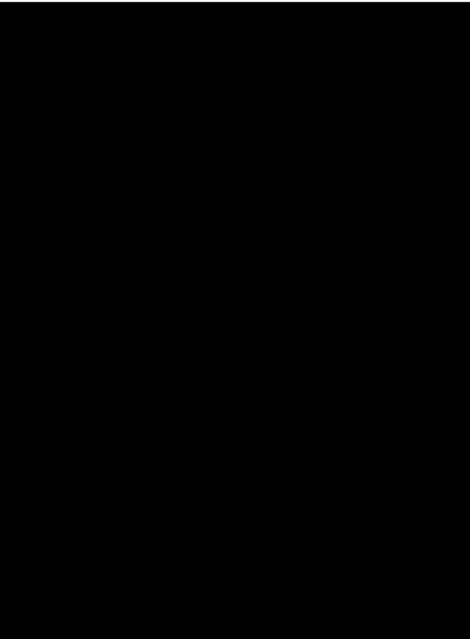
DRAM

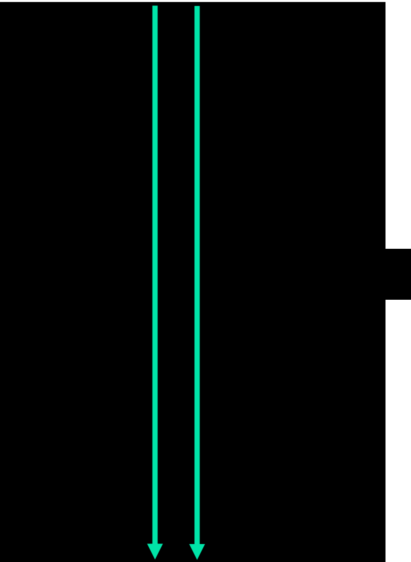
- Smaller than SRAM
- Require data refresh to compensate for leakage

1-Transistor DRAM Cell (preclass 3)

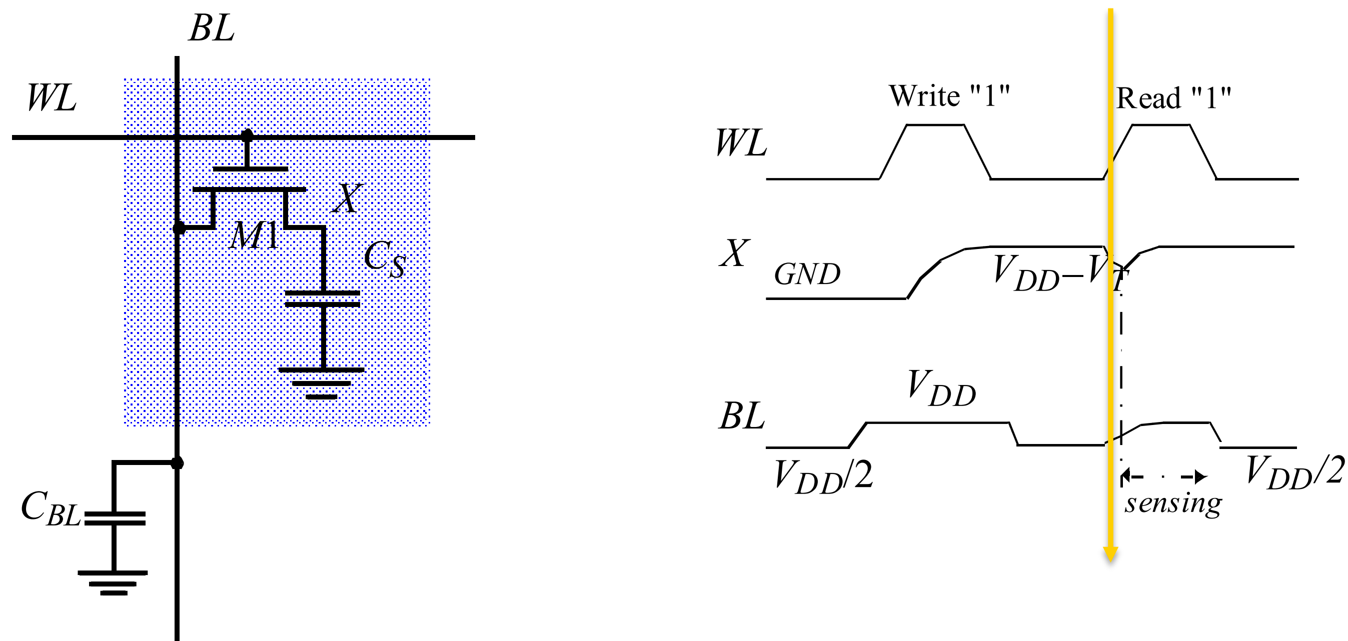


1-Transistor DRAM Cell





1-Transistor DRAM Cell

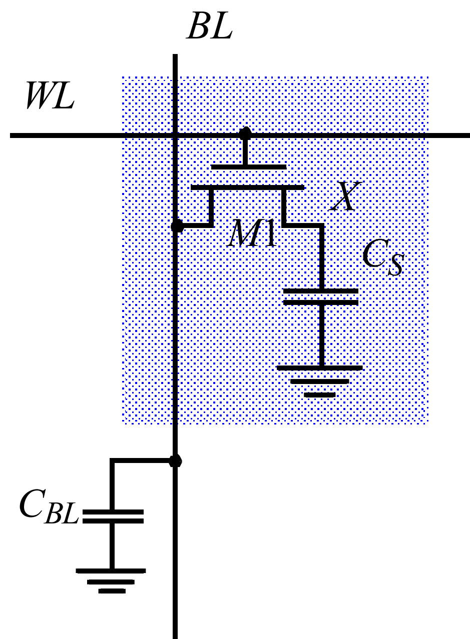


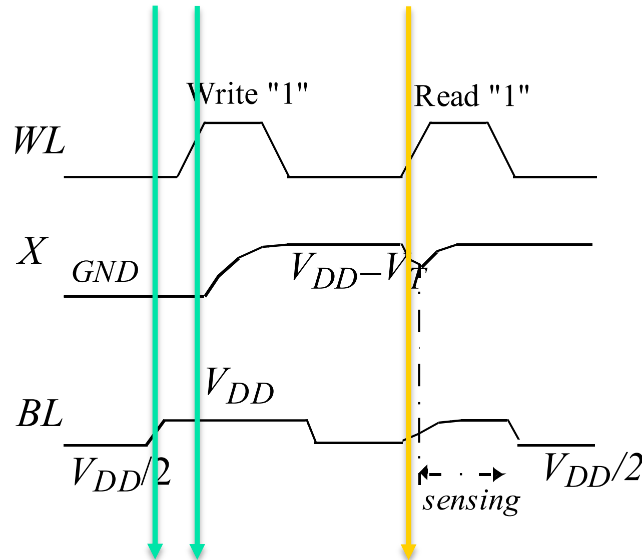
Write: C_S is charged or discharged by asserting WL and BL.
Read: Charge redistribution takes places between bit line and storage capacitance

$$\Delta V = V_{BL} - V_{PRE} = (V_{BIT} - V_{PRE})\frac{C_S}{C_S + C_{BL}}$$

Voltage swing is small; typically around 250 mV.

1-Transistor DRAM Cell





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Voltage swing is small; typically around 250 mV.

DRAM Cell Observations

- 1T DRAM requires a sense amplifier for each bit line, due to charge redistribution read-out
- DRAM memory cells are single ended in contrast to SRAM cells
- The read-out of the 1T DRAM cell is destructive; read and refresh operation are necessary for correct operation
- Unlike 3T cell, 1T cell requires presence of an extra capacitance that must be explicitly included in the design
- When writing a "1" into a DRAM cell, a threshold voltage is lost. This loss can be circumvented by bootstrapping the word lines to a higher value than V_{DD}.

Big Idea

- 6T SRAM
  - Robust cell when sized carefully
- 5T SRAM
  - More sensitive to sizing than 6T SRAM
- Minimize area of repeated cell
  - 6T/5T SRAM
  - Mulitport trade off area for function
  - 1T/3T DRAM helps but slower
- Compensate with periphery
  - Decoders
  - Bitline (column) drivers
  - Sensing/Amplification (regeneration/restoration)

Admin

- Project 2 out – START ASAP!
  - Work in teams up to 2
  - Final report due 4/30
- Wednesday 4/16 Midterm 2 (next week)
  - 1:45pm-3:45pm **in class**
  - **Midterm 2 Review session (4/16) - in class**
  - Lectures 11-18
  - Closed note, calculator allowed
  - All old exams online
    - 2015-2024
  - **Do review your preclass!!**

Acknowledgement

- Prof. André DeHon (University of Pennsylvania)
- Prof. Tania Khanna (University of Pennsylvania)

Additional Reading (Optional)

Dual-Ported SRAM

- Simple dual-ported SRAM
  - Two independent single-ended reads
  - Or one differential write

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  - Two independent single-ended reads
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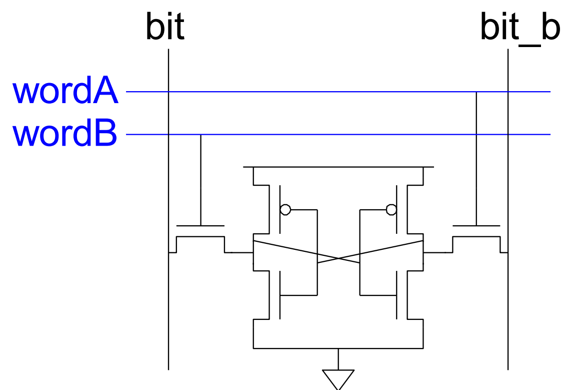
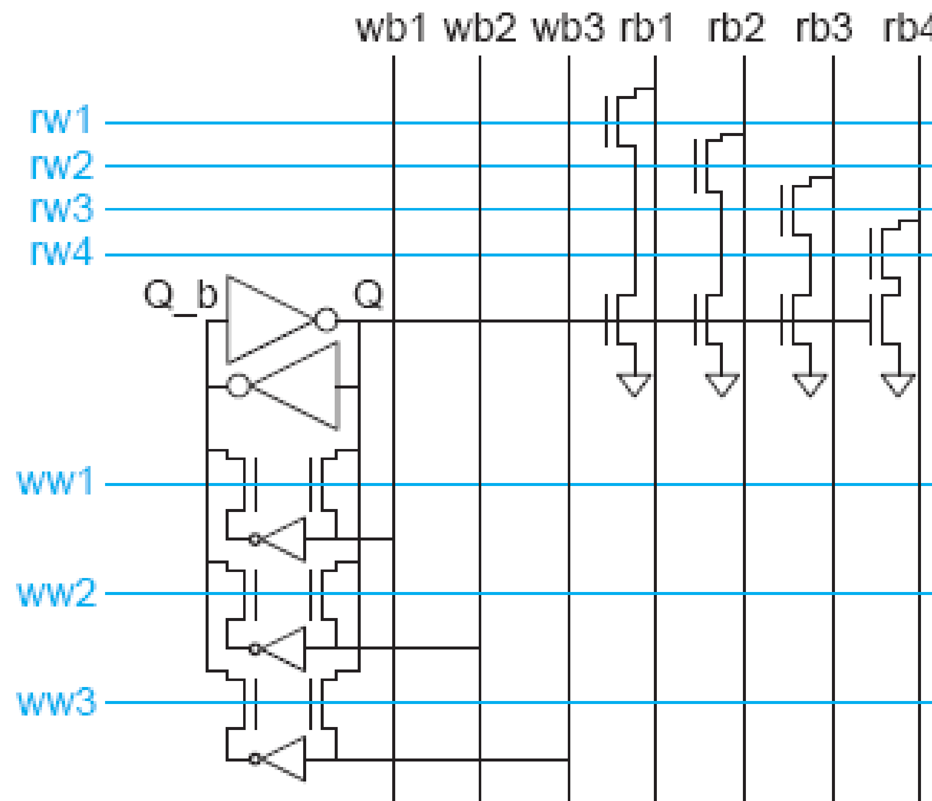


- Do two reads and one write by time multiplexing
  - Read during ph1, write during ph2

Multi-Ported SRAM

- Adding more access transistors hurts read stability
- Multiported SRAM isolates reads from state node
- Single-ended bitlines save area



Register File Cell

- Single-ended 2-read/1-write ports (Slow-write)

bitWR
bitRD1
bitRD2
I1
I3
A
N2
N1
N3
N4
I2
wordWR
wordRD1
wordRD2
wrEN

Register File Cell

- Single-ended Read/Dual-ended Write